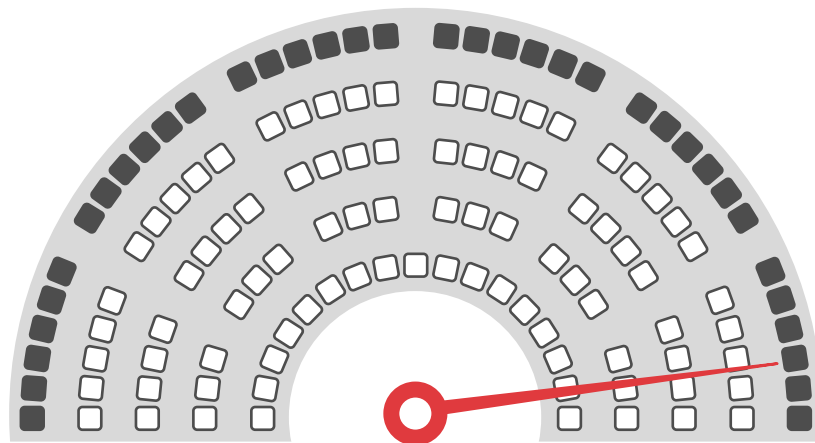




Optimized Democracy

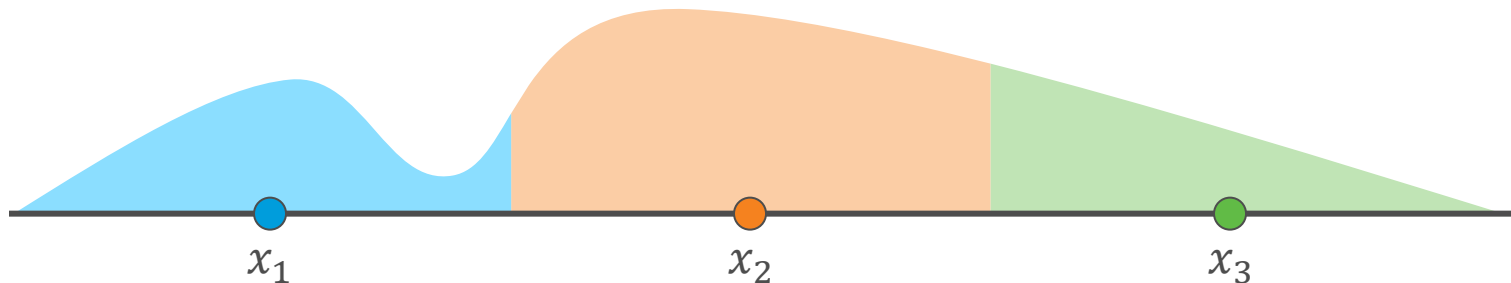
Fall 2026 | Lecture 7

Electoral Competition

Ariel Procaccia | Harvard University

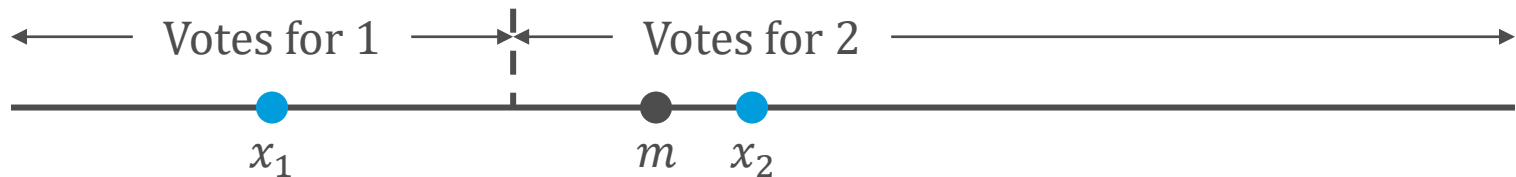
THE HOTELLING MODEL

- Political spectrum is $\mathbb{R}$
- There is a continuous distribution F of voters, each with a peak in $\mathbb{R}$
- Players are candidates, who strategically choose positions $x_1, \dots, x_n$
- Each candidate attracts the votes of voters who are closest to them, with votes being split equally in case of a tie



THE HOTELLING MODEL

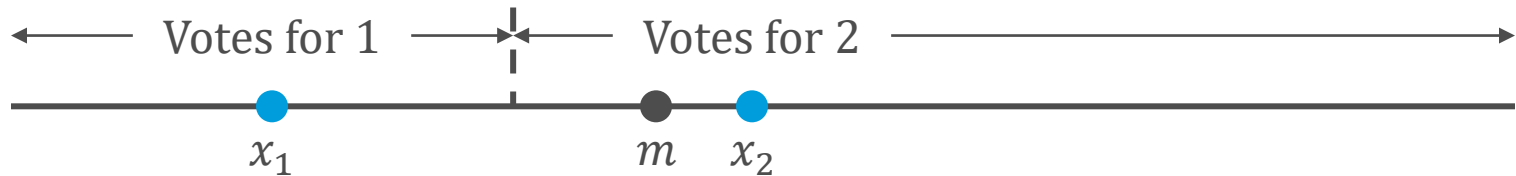
- Two candidates seek to win a plurality of votes
- The utility of each candidate is 1 if they win, $1/2$ if they tie, and 0 if they lose
- Denote the median peak by m (assume for simplicity that it's unique)



Who wins?

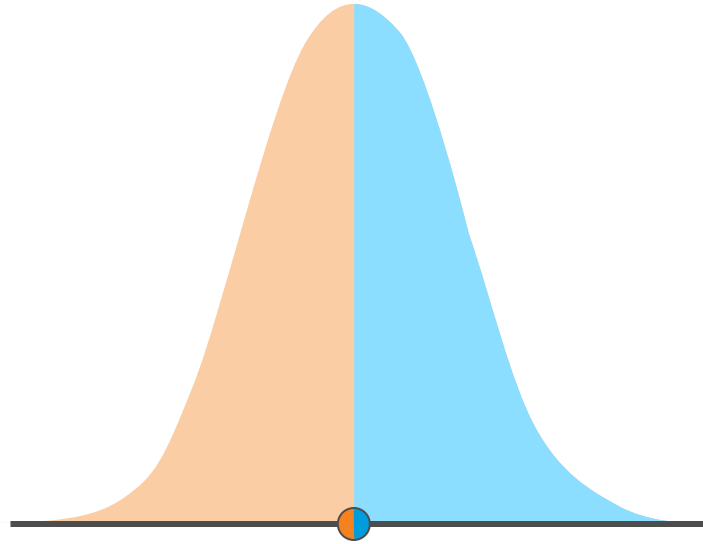
NASH EQUILIBRIUM

A **Nash equilibrium** is a profile $\mathbf{x} \in \mathbb{R}^n$ such that each player is **best responding** to the others, i.e., for each player i and alternative strategy $x'_i \in \mathbb{R}$, $u_i(\mathbf{x}) \geq u_i(x'_i, \mathbf{x}_{-i})$



Is this a Nash equilibrium?

THE MEDIAN VOTER THEOREM



Theorem: In the Hotelling Model with two candidates, there is always a unique Nash equilibrium at (m, m)

PROOF OF THEOREM

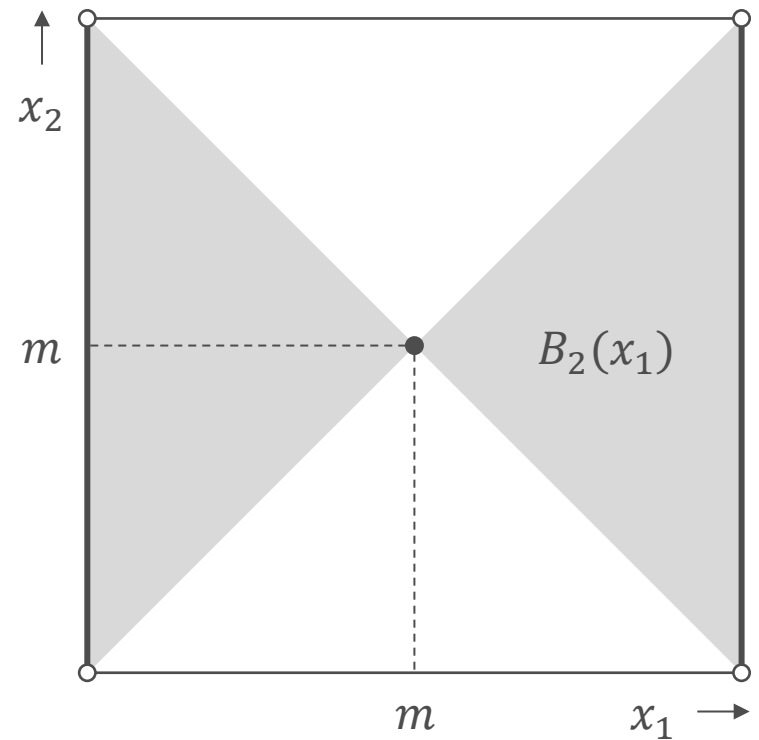
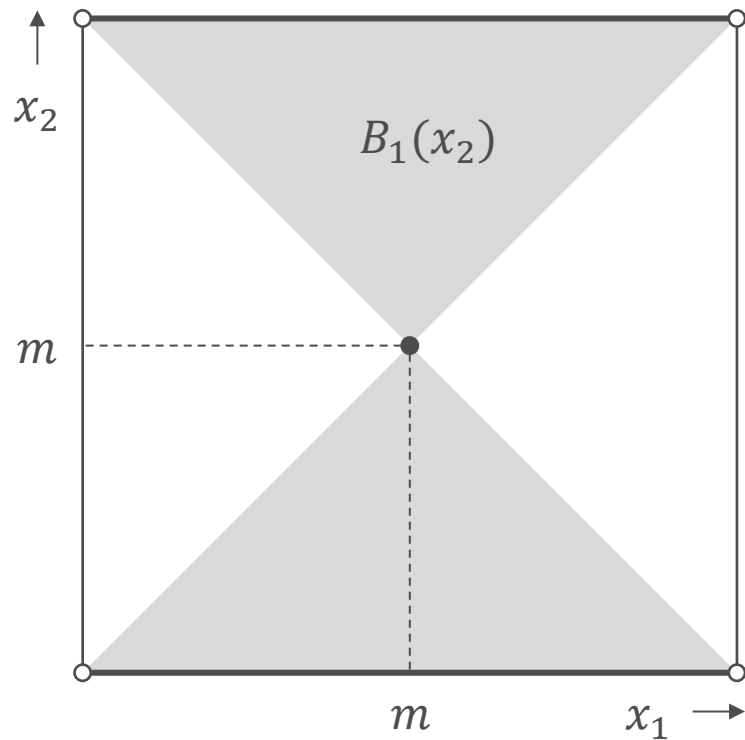
- If $x_2 < m$, the best response for 1, denoted $B_1(x_2)$, is all positions x_1 such that

$$x_1 > x_2 \quad \text{and} \quad \frac{x_1 + x_2}{2} < m$$

- A symmetric argument holds if $x_2 > m$
- If $x_2 = m$, the best response for 1 is m
- Therefore, it holds that

$$B_1(x_2) = \begin{cases} \{x_1 : x_2 < x_1 < 2m - x_2\} & x_2 < m \\ \{m\} & x_2 = m \\ \{x_1 : 2m - x_2 < x_1 < x_2\} & x_2 > m \end{cases}$$

PROOF OF THEOREM



The unique Nash equilibrium is at (m, m)

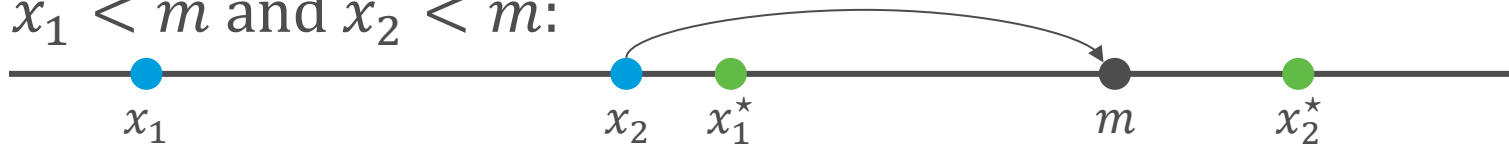
POLICY-MOTIVATED CANDIDATES

- What if candidates care about policy and not just about winning?
- Suppose i has a preferred position x_i^* , and their utility depends on the distance between x_i^* and the position of the winner
- If there's a tie then candidates evaluate the induced lottery over winning positions
- **Theorem:** If $x_1^* < m < x_2^*$ then (m, m) is the unique Nash equilibrium

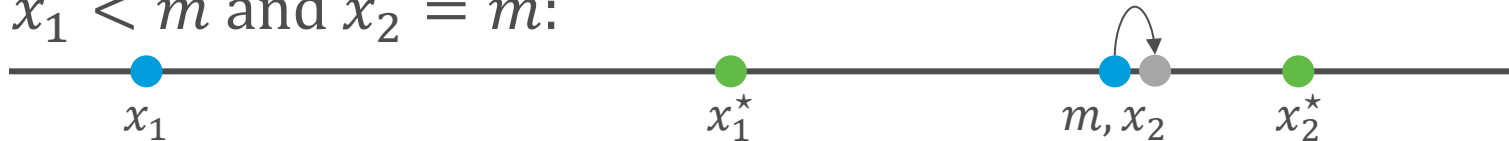
PROOF SKETCH

Rule out cases for which $(x_1, x_2) \neq (m, m)$:

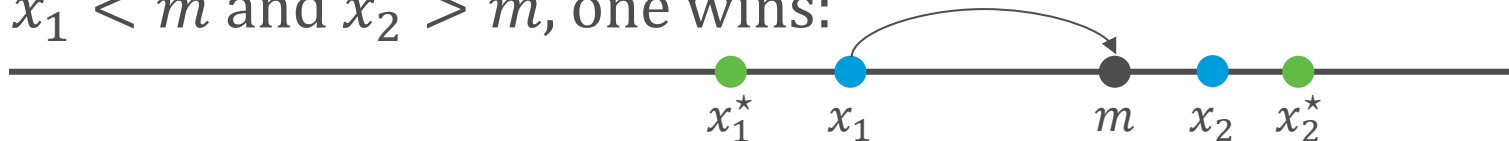
$x_1 < m$ and $x_2 < m$:



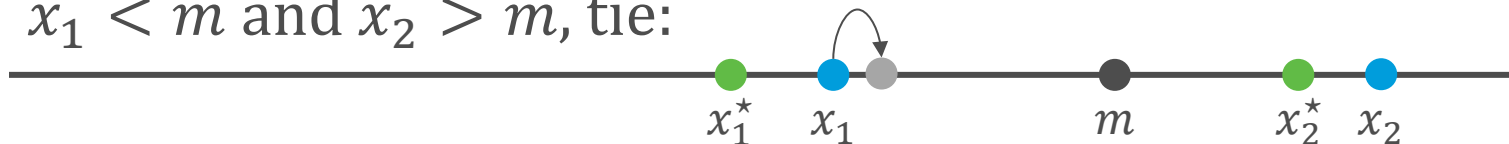
$x_1 < m$ and $x_2 = m$:



$x_1 < m$ and $x_2 > m$, one wins:



$x_1 < m$ and $x_2 > m$, tie:





Harold Hotelling

1895–1973

“The competition for votes between the Republican and Democratic parties does not lead to a clear drawing of issues, an adoption of two strongly contrasted positions between which the voter may choose. Instead, each party strives to make its platform as much like the other's as possible.”





“The competition for votes between the Republican and Democratic parties does not lead to a clear drawing of issues, an adoption of two strongly contrasted positions between which the voter may choose. Instead, each party strives to make its platform as much like the other’s as possible.”

(m, m) USED TO MAKE SENSE

TOWARD A MORE RESPONSIBLE TWO-PARTY SYSTEM

A Report of the Committee on Political Parties
American Political Science Association

Vol. XLIV

September, 1950

Number 3, Part 2

CITIZEN CANDIDATES

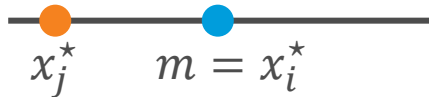
- What if citizens can run in the election?
- Continuum of voters as before, which are now the players
- The position of player i is $x_i^\star$
- Each player chooses whether to run or not
- Players who run are “honest” about their position
- The cost of running is c , and the benefit of winning is w

CITIZEN CANDIDATES

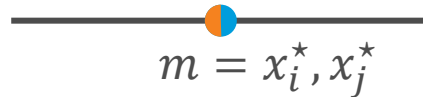
- The utility of player i is
 - $-|x_i^* - x_j^*|$ if i doesn't run and j wins
 - $-|x_i^* - x_j^*| - c$ if i runs and j wins
 - $w - c$ if i runs and wins
 - $-\infty$ if nobody runs (for simplicity)
- In case of a tie, each player gets their expected utility
- **Theorem:** There is a one-candidate equilibrium if and only if $w \leq 2c$, where if $c \leq w \leq 2c$ then the candidate's position is m , and if $w < c$ then it is in $\left[m - \frac{c-w}{2}, m + \frac{c-w}{2}\right]$

PROOF OF THEOREM

- If $w \leq 2c$, there is an equilibrium where a single player with position m runs:



If another candidate j with a different position runs, they lose



If another candidate j at m runs, they get $w/2 - c \leq 0$ instead of 0

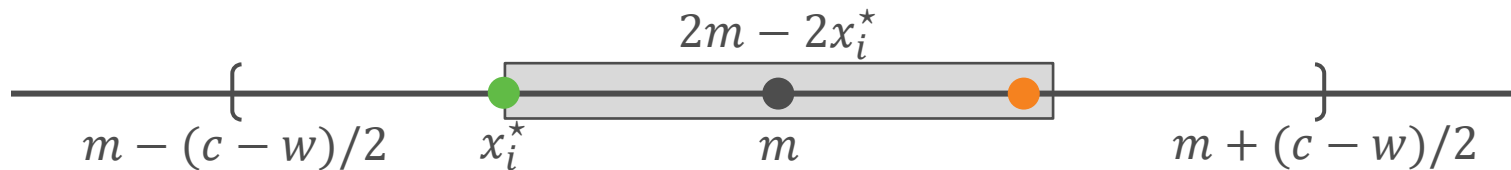


If the single candidate i drops out, they get $-\infty$

- If $w > 2c$, for any profile where one player enters, another player with the same position would wish to enter

PROOF OF THEOREM

- Consider a single candidate with w.l.o.g. $x_i^* < m$
- Any candidate j in the interval $(x_i^*, 2m - x_i^*)$ can run, win, and “pay” $c - w$ instead of $x_j^* - x_i^*$
- Since $x_j^* - x_i^*$ can be arbitrarily close to $2m - 2x_i^*$, for this to be an equilibrium it must hold that $c - w \geq 2m - 2x_i^*$ (and hence $c > w$)
- These conditions are also sufficient ■



CITIZEN CANDIDATES

- We next (briefly) consider equilibria with two candidates

Poll 1

Is there a two-candidate equilibrium with both candidates at m ?

- Always
- Sometimes
- Never



EXTENSIONS

- We introduced policy-motivation and citizen candidates into the original Hotelling model, but there are other gaps from reality

Poll 2

What are some other ingredients that are missing from the model?



BIBLIOGRAPHY

M. J. Osborne. **Spatial Models of Political Competition Under the Plurality Rule: A Survey of Some Explanations of the Number of Candidates and the Positions They Take.** Canadian Journal of Economics, 1995.

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