

Optimized Democracy

Spring 2021, Lecture 9, 2021-02-24

Participatory Budgeting

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Outline

- Aside on Single Transferable Vote
- More on approval-based committee elections: Phragmén's rule
- Participatory Budgeting
- *Repeating theme*: Can get proportionality by explicitly dividing “voting power” equally among voters. (Rather than magically proportional PAV.)

Single Transferable Vote

6	4	5	2
<i>a</i>	<i>b</i>	<i>c</i>	<i>b</i>
<i>b</i>	<i>c</i>	<i>a</i>	<i>a</i>
<i>c</i>	<i>a</i>	<i>b</i>	<i>c</i>

6	4	5	2
<i>a</i>	<i>b</i>		<i>b</i>
<i>b</i>		<i>a</i>	<i>a</i>
	<i>a</i>	<i>b</i>	

6	4	5	2
<i>a</i>			
		<i>a</i>	<i>a</i>
	<i>a</i>		

6	4	5	2
<i>a</i>	<i>b</i>	<i>c</i>	<i>a</i>
<i>b</i>	<i>c</i>	<i>a</i>	<i>b</i>
<i>c</i>	<i>a</i>	<i>b</i>	<i>c</i>

6	4	5	2
<i>a</i>		<i>c</i>	<i>a</i>
	<i>c</i>	<i>a</i>	
<i>c</i>	<i>a</i>		<i>c</i>

6	4	5	2
		<i>c</i>	
	<i>c</i>		
<i>c</i>			<i>c</i>

Single Transferable Vote for Committees

- STV can also be used to elect a k -committee.
- Initially, each voter gets a ‘budget’ of \$1.
- It costs $\$n/k$ to elect a candidate.
- As long as there is a candidate that is ranked first by voters who together have at least $\$n/k$, elect the candidate and charge those voters $\$n/k$.
- Otherwise, eliminate the candidate whose supporters are poorest, and repeat.
- Exercise: Show STV elects k candidates.

Proportionality for Solid Coalitions (PSC)

- Suppose a set $S \subseteq N$ with $|S| \geq \ell \frac{n}{k}$ has the same set $T \subseteq C$ of $|T| = t \geq \ell$ candidates they all rank top, so $T \succ C \setminus T$ for all $i \in S$ (not necessarily ranked in the same order).
- Then $|W \cap T| \geq \ell$.
- STV satisfies this! (no matter how spending is distributed)

STV satisfies PSC

- Let $S \subseteq N$ with $|S| \geq \ell \frac{n}{k}$ agree on $|T| = t \geq \ell$ candidates.
- Suppose PSC failed for S . Then there is a time when
 - $\ell - j$ candidates from T have been elected
 - j further candidates from T need to be elected for PSC
 - all but j candidates from T have been elected or eliminated.
- Group S has only paid at most $(\ell - j) \frac{n}{k}$ thus far, so has $j \frac{n}{k}$ left over. So at least one of the j candidates has $\frac{n}{k}$ support, and this will remain true until all j candidates have been elected.

Hare vs Droop Quota

- The value “ $\frac{n}{k}$ ” is known as the **Hare quota**.
 - Intuition: electorate is split into equal-sized groups, each of which is assigned one seat.
- But we can also use $\frac{n}{k+1} + \varepsilon$, the **Droop quota**.
 - This works because there are at most k disjoint subsets of N of size $\frac{n}{k+1} + \varepsilon$.
 - Guarantees representation to smaller groups.
 - For $k = 1$, this says majority needs to be followed.
- Everything we’ve said works for Droop quota if we are more careful in the proofs.
 - PAV satisfies Droop EJR, Droop-STV satisfies Droop PSC

Open Problem

Does there exist a ranking-based committee rule that is monotonic and satisfies PSC?

Recap: Approval-based Committee Elections

- Proportional Approval Voting maximizes $\sum_i 1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{|W \cap A_i|}$.
- PAV satisfies Extended Justified Representation:
If $S \subseteq N$ with $|S| \geq \ell \frac{n}{k}$ agrees on ℓ candidates $T \subseteq \bigcap_{i \in S} A_i$, then $|W \cap A_i| \geq \ell$ for some $i \in S$.
- PAV is NP-complete to compute.
- Sequential PAV fails EJRR even for $\ell = 1$.
- *Question:* Can we get something proportional in polynomial time?

Is PAV always right?

$$k = 12$$

4	5	6	10	14	18
3			9	13	17
2			8	12	16
1			7	11	15
v_1	v_2	v_3	v_4	v_5	v_6

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EJR not strong enough to capture this!

Phragmén's Rule

- Proposed in 1894. Thiele proposed PAV in 1895. Phragmén criticized it in 1899, for a reason similar to

4	5	6	10	14	18
3			9	13	17
2			8	12	16
1			7	11	15



EDVARD PHRAGMÉN

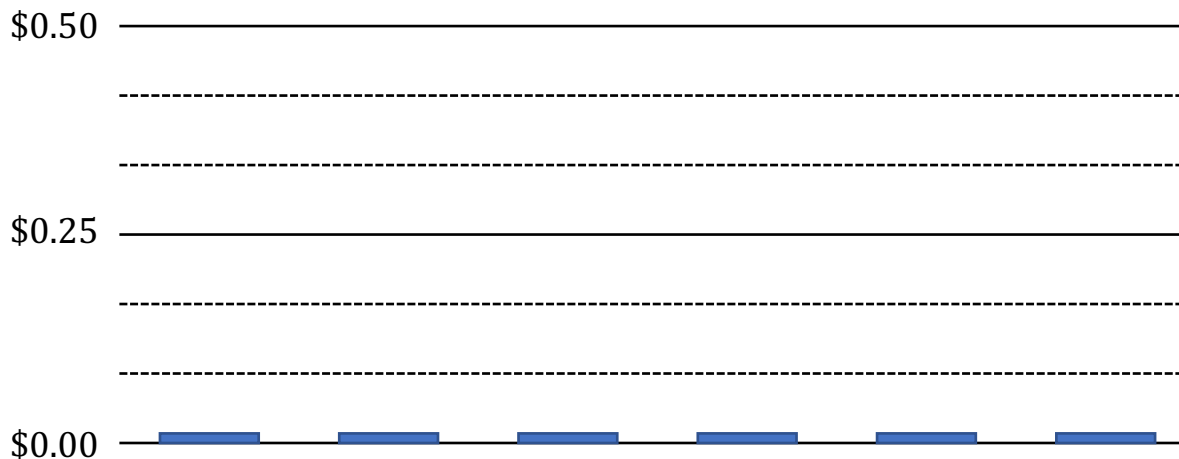
- Phragmén's proposal:
 - Each voter starts with a bank account with \$0.
 - Fill bank accounts at the same rate, until the approvers of some unelected candidate together hold $\$ \frac{n}{k}$.
 - Elect the candidate and reset approvers' accounts to \$0.
 - Stop after k candidates are elected.

Phragmén's Rule: Example

$k = 12$



Sveriges Riksbank



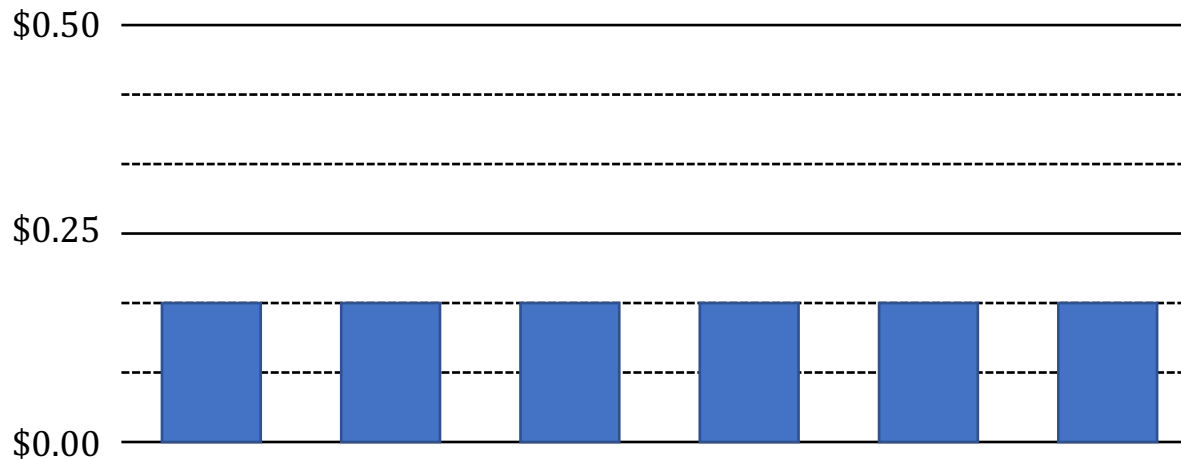
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Sveriges Riksbank



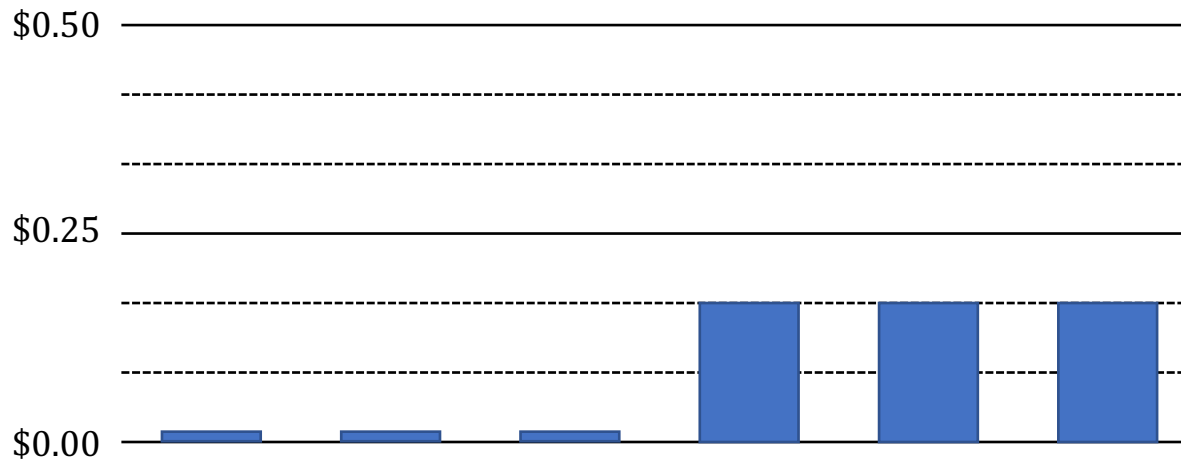
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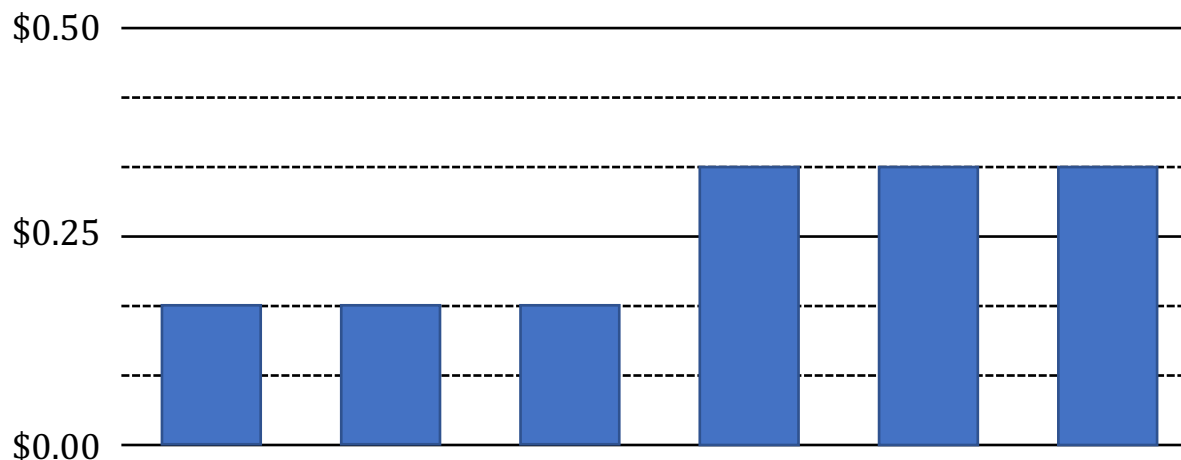
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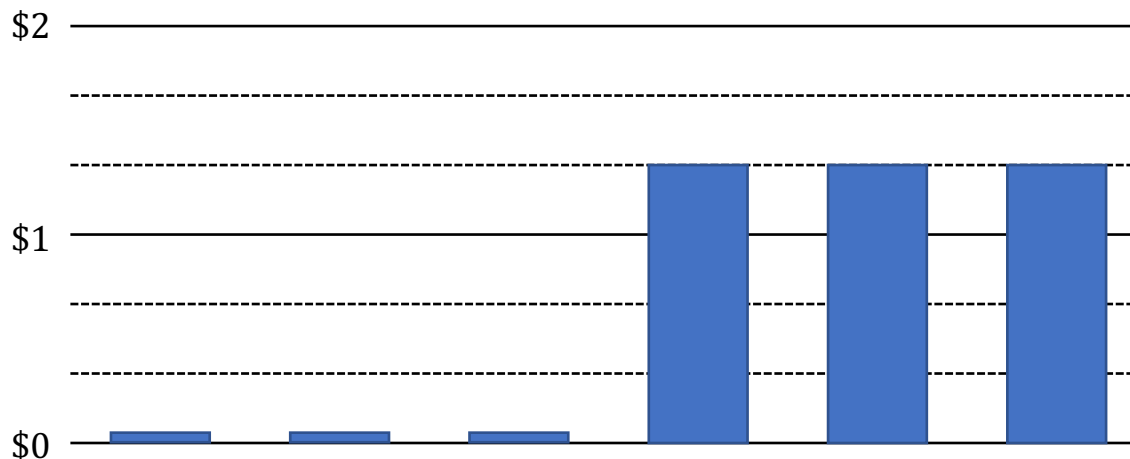
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Phragmén's Rule: Example

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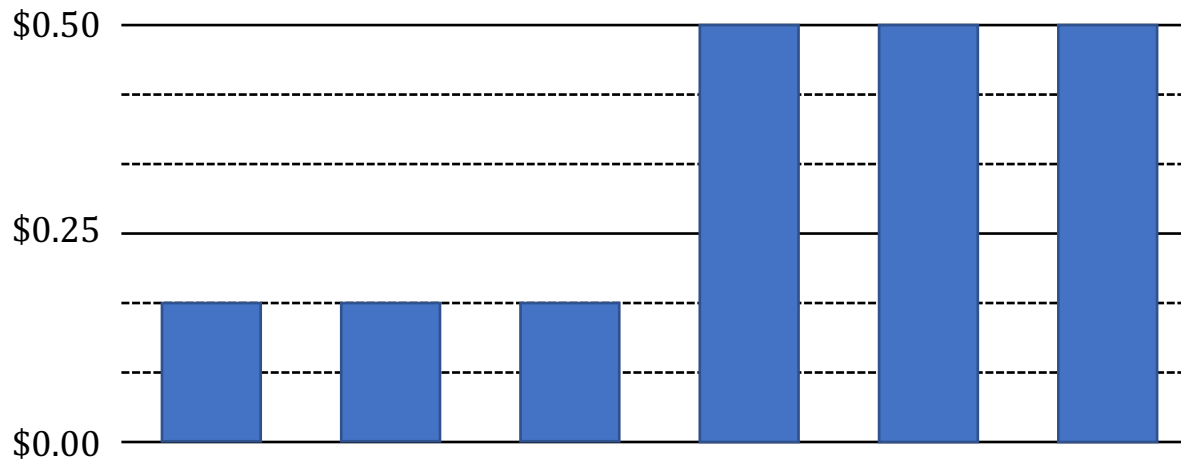
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Phragmén's Rule: Example

$k = 12$



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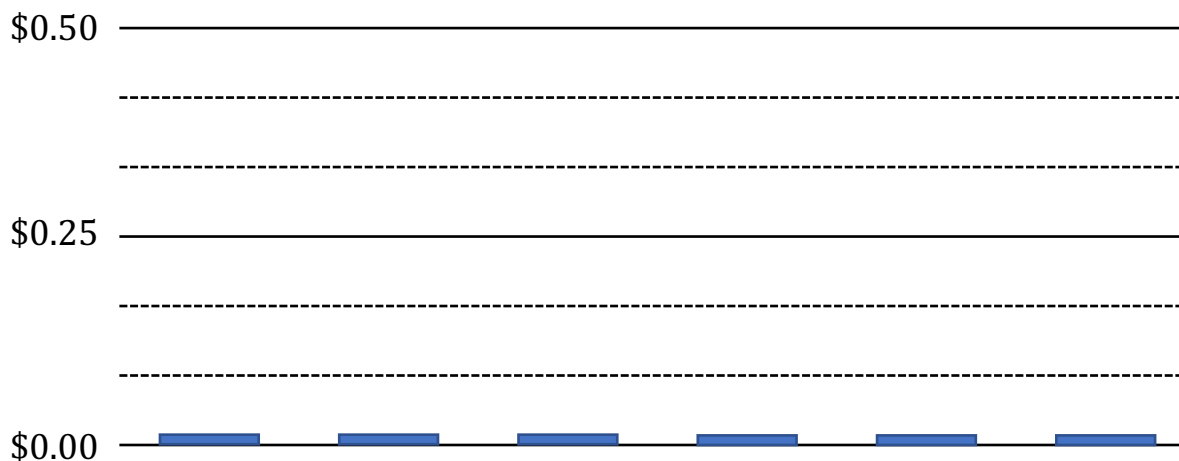
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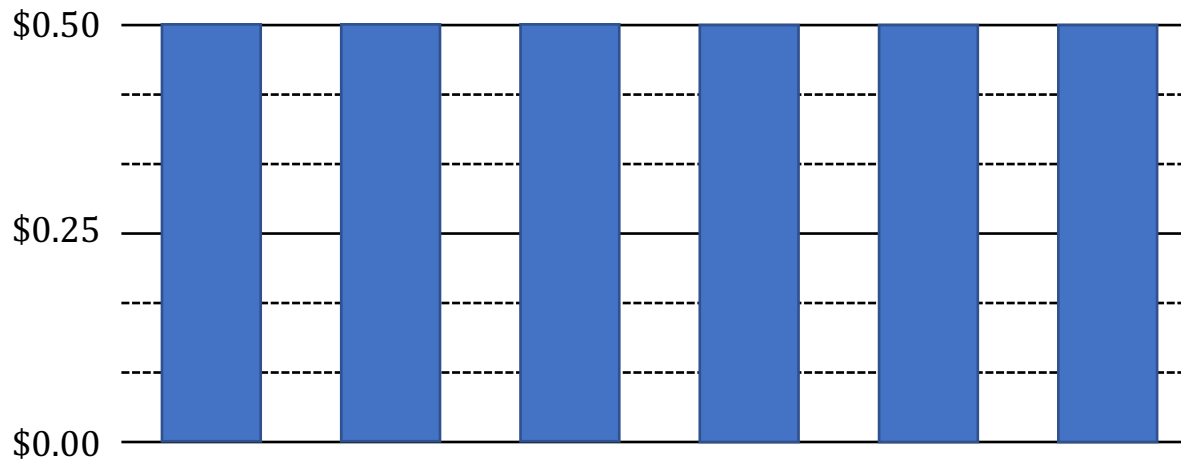
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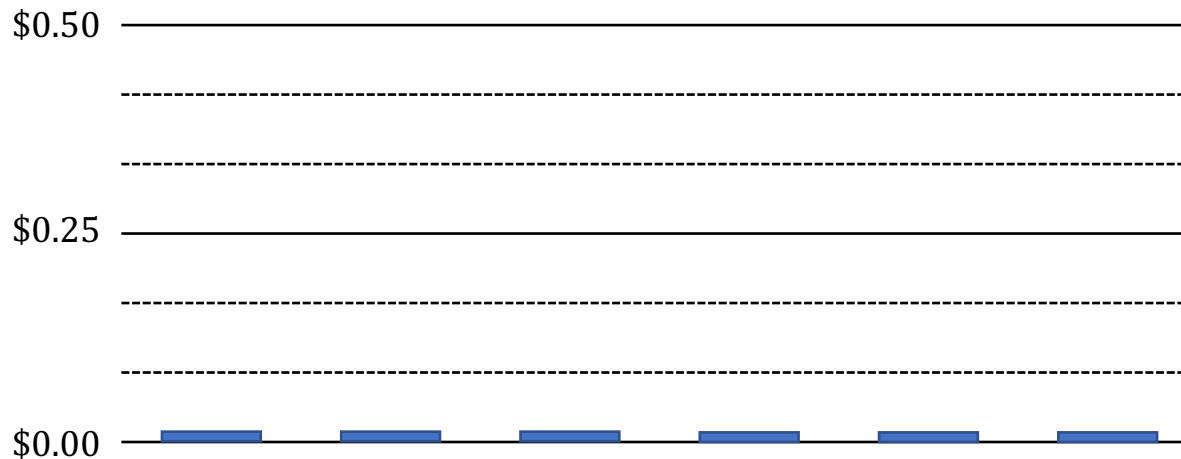
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Phragmén's Rule: Example

$$k = 12$$



Sveriges Riksbank



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v_1	v_2	v_3	v_4	v_5	v_6

Phragmén's Rule: Proportionality

- Phragmén's rule violates EJR (largish example with 24 voters, 14 candidates, $k = 12$).

- But it satisfies a weaker version (“**PJR**”):

If $S \subseteq N$ with $|S| \geq \ell \frac{n}{k}$ agrees on ℓ candidates
 $T \subseteq \bigcap_{i \in S} A_i$, then $|W \cap \bigcup_{i \in S} A_i| \geq \ell$.

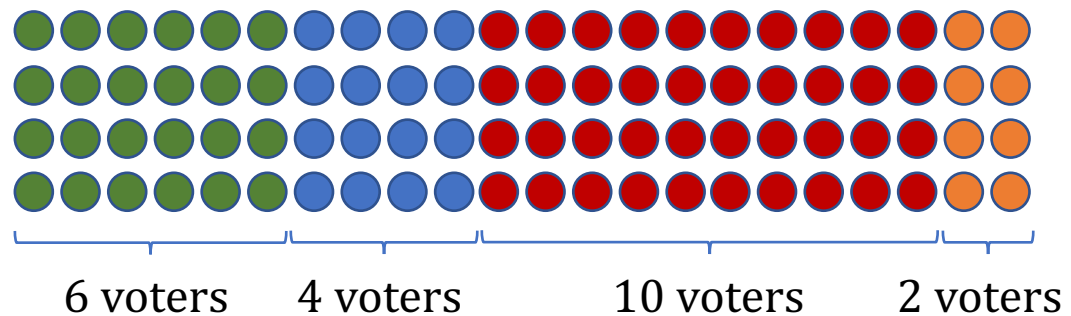
6	7	8	9	10
5				
4				
3				
2				
1				
v_1	v_2	v_3	v_4	v_5

Phragmén's Rule: PJR

- If $S \subseteq N$ with $|S| \geq \ell \frac{n}{k}$ agrees on ℓ candidates $T \subseteq \bigcap_{i \in S} A_i$, then $|W \cap \bigcup_{i \in S} A_i| \geq \ell$.
- *Proof:* By the time Phragmén terminates, each voter has received at least \$1.
- If it terminates exactly at the \$1-point, then all money was spent. So S spent $\ell \frac{n}{k}$, and so they bought ℓ candidates from $\bigcup_{i \in S} A_i$.
- If it terminates strictly later, consider \$1-point. If then $|W \cap \bigcup_{i \in S} A_i| \leq \ell - 1$, then S now holds at least $\frac{n}{k}$, so can purchase a candidate from T .

Proportional Rankings

- Note: you don't have to stop Phragmén after it has elected k candidates (same for SeqPAV)
- This way, we get a *proportional ranking*.
- In particular, every prefix satisfies PJR. (Or think of party-list profiles.)
- Applications:
 - Ranking comments by upvotes
 - Displaying proposal variants in LiquidFeedback
- **Open Problem:** Do there exist EJR rankings?



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BUDGET PARTICIPATIF

VOS PROJETS PRÉFÉRÉS
DU 6 AU 22 SEPTEMBRE 2019

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ÉTOILE

ÉTOILE



You have selected **1 / 4** projects.

BCYF Hyde Park Dance Studio Renovation


A renovated dance studio at the Hyde Park Community Center for children of all ages.

Estimated Cost: \$286,000

Location: BCYF Hyde Park Community Center, Hyde Park

Select



 Click image for slide show

Bike Lane Installation

After a study, bike lanes will begin to be installed around Charlestown Navy Yard, Bunker Hill housing, and Charlestown High.


Estimated Cost: \$200,000

Location: Charlestown

✓ Selected

Remove



 Click image for slide show

Wicked Free Wifi 2.0


Wicked free Wi-Fi 2.0 provides Wi-fi at locations with young people.

Estimated Cost: \$119,000

Location: Various High Schools and Community Centers, Dorchester, Roxbury, East Boston, Charlestown

Select



 Click image for slide show



You still have **\$1,350,000** left.



Education

Bathroom Renovations at M.S./H.S. 223

Renovation of girls' and boys' bathrooms including stalls, lighting, painting, and having walls re-glazed.

Estimated Cost: \$150,000

Location: 360 E.145th St. (Bronx - Mott Haven)

✓ Selected

Remove

Technology Upgrades

Technology upgrades for Park East High School and Central Park East High School.

Estimated Cost: \$312,000

Location: 230 E.105th; 1573 Madison Ave. (El Barrio/East Harlem)

Select

Air Conditioning at Bronx Schools

Installation of 1 air conditioning system at 345 Brooke Avenue for schools X343, X224 & X334. Installation of 1 air conditioning system at PS 161x.

Estimated Cost: \$500,000

Location: 628 Tinton Ave, 345 Brook Ave. (Bronx - Mott Haven & Longwood)



Air Conditioning: P.S.179, P.S.369, P352

Installation of two HVAC units at P.S. 179X, P.S. 369X and P.S. 352X.

Estimated Cost: \$500,000

Location: 468 E. 140 St. (Bronx - Mott Haven)

✓ Selected

Remove

4e arrondissement

Greedy: total utility **3 500**. Funds 5 projects, avg cost 293 000 🚧🚧🧑🌳💡

Optimal: total utility **6 878**. Funds 14 projects, avg cost 98 928 🚧🚧🏛️🏛️🏛️🧑🧑🧑💡💡🔧🔧🔧🚲

Project Name	QPOP	Cost €	Votes	v / k€	Greedy	Optimal
🚧 Un mur végétalisé au croisement des rues Blancs Manteaux et Archiv		30 000	788	26	✓	✓
💡 Un café solidaire dans le quartier de la tour Saint-Jacques		15 000	706	47	✓	✓
🧑 Une salle d'arts plastiques pour l'école Saint-Merri Renard		300 000	702	2	✓	✓
🌳 Rénovation énergétique exemplaire d'une école du 4e		1 000 000	655	1	✓	–
🚧 Végétalisation de la rue de l'Arsenal		120 000	649	5	✓	✓
🧑 Un collège Charlemagne accessible aux personnes à mobilité réduit		200 000	630	3	–	✓
🔧 Faire du hall d'accueil de la piscine Saint-Merri un lieu de convivia		20 000	528	26	–	✓
🔧 Des agrès sportifs place des Vosges		15 000	491	33	–	✓
🏛️ Mise en valeur des pierres de la prison de la Bastille		20 000	473	24	–	✓
💡 Un fauteuil roulant électrique pour rompre l'isolement		5 000	453	91	–	✓
🔧 Création d'un auvent sur une aire de jeux d'un square du 4e		150 000	410	3	–	✓
🧑 L'Ascenseur, un tiers-lieu pour l'égalité des chances ouvert sur le		350 000	315	1	–	✓
🏛️ Baliser le passage du chemin de Compostelle dans le 4e		30 000	265	9	–	✓
🚲 Des arceaux vélos rue de la Reynie		10 000	240	24	–	✓
🏛️ Faciliter l'accès au cellier d'Ourscamp		120 000	228	2	–	✓
				Total €	1 465 000	1 385 000

Participatory Budgeting: Model

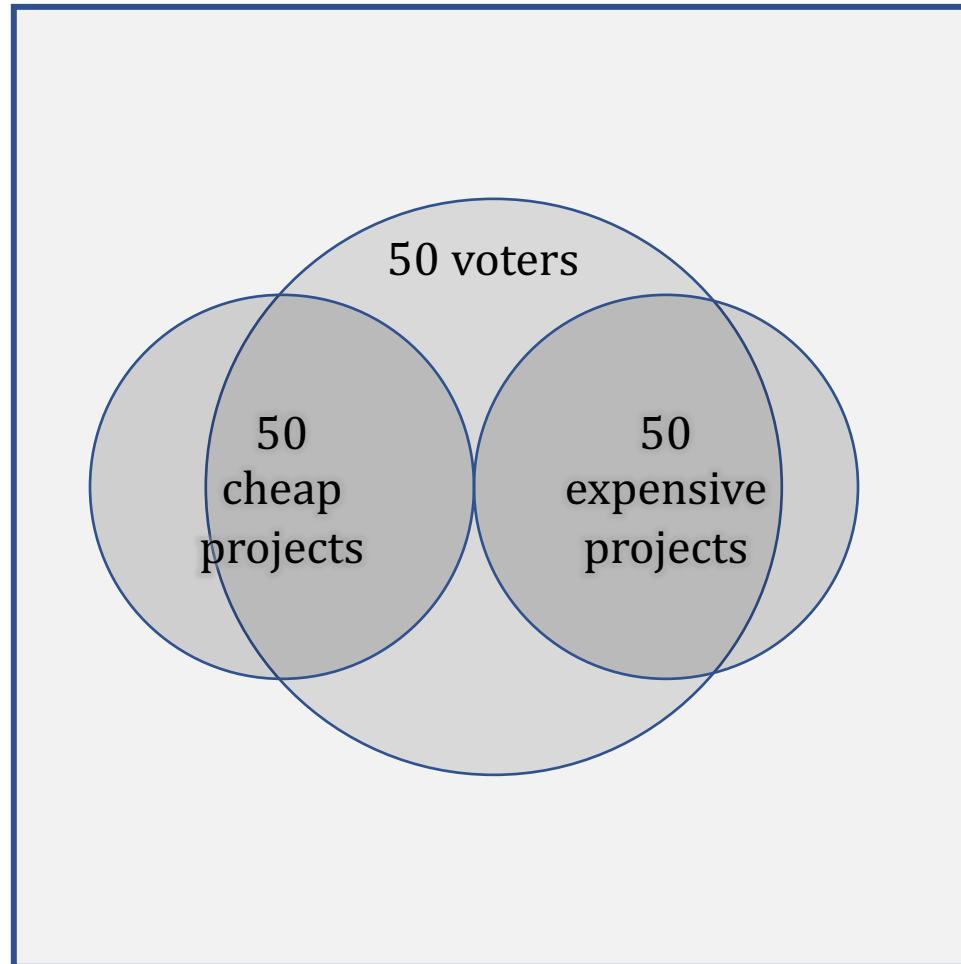
- A set \mathcal{C} of projects, each with a cost
- A budget limit B
- Outcome: set $W \subseteq \mathcal{C}$ with $\sum_{c \in W} \text{cost}(c) \leq B$.
- A set N of n voters
- Each voter $i \in N$ approves a subset $A_i \subseteq \mathcal{C}$.
- Mostly, we say that i 's utility is $u_i(W) = |A_i \cap W|$ (this is a dichotomous preference assumption).
- **Unit cost assumption:** $\text{cost}(c) = 1$ for all c .

Three interpretations of “AV”

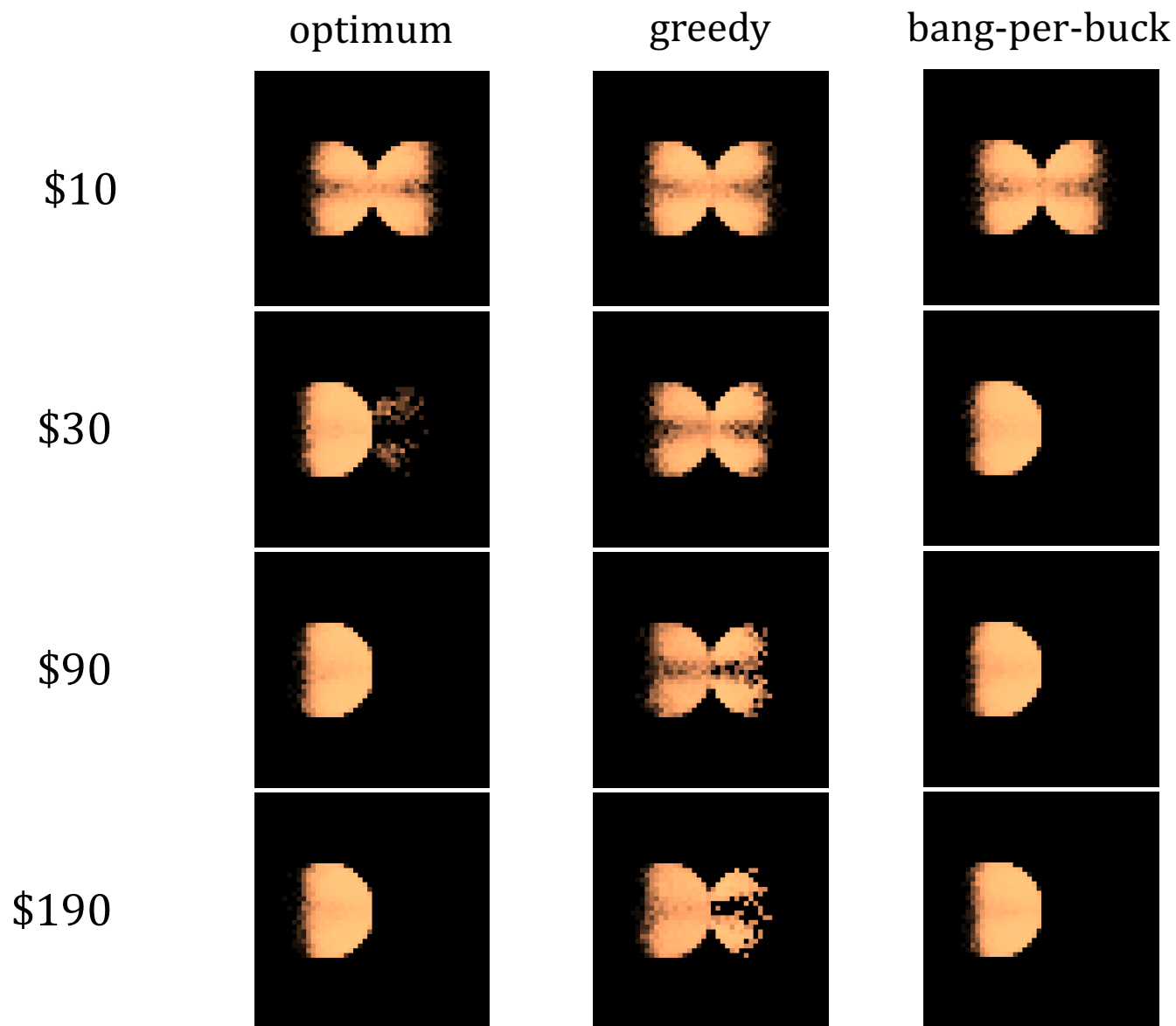
- Optimize $\sum_{i \in N} u_i(W) = \sum_{c \in W} \text{approval-score}(c)$.
- Greedy: add projects in order of approval score, skipping unaffordable projects.
- Bang-per-buck greedy: add projects in order of approval score divided by cost.

Experiments

Budget = \$1000. Cheap = \$10. Expensive = \$10, \$30, \$90, \$190.



Experiments

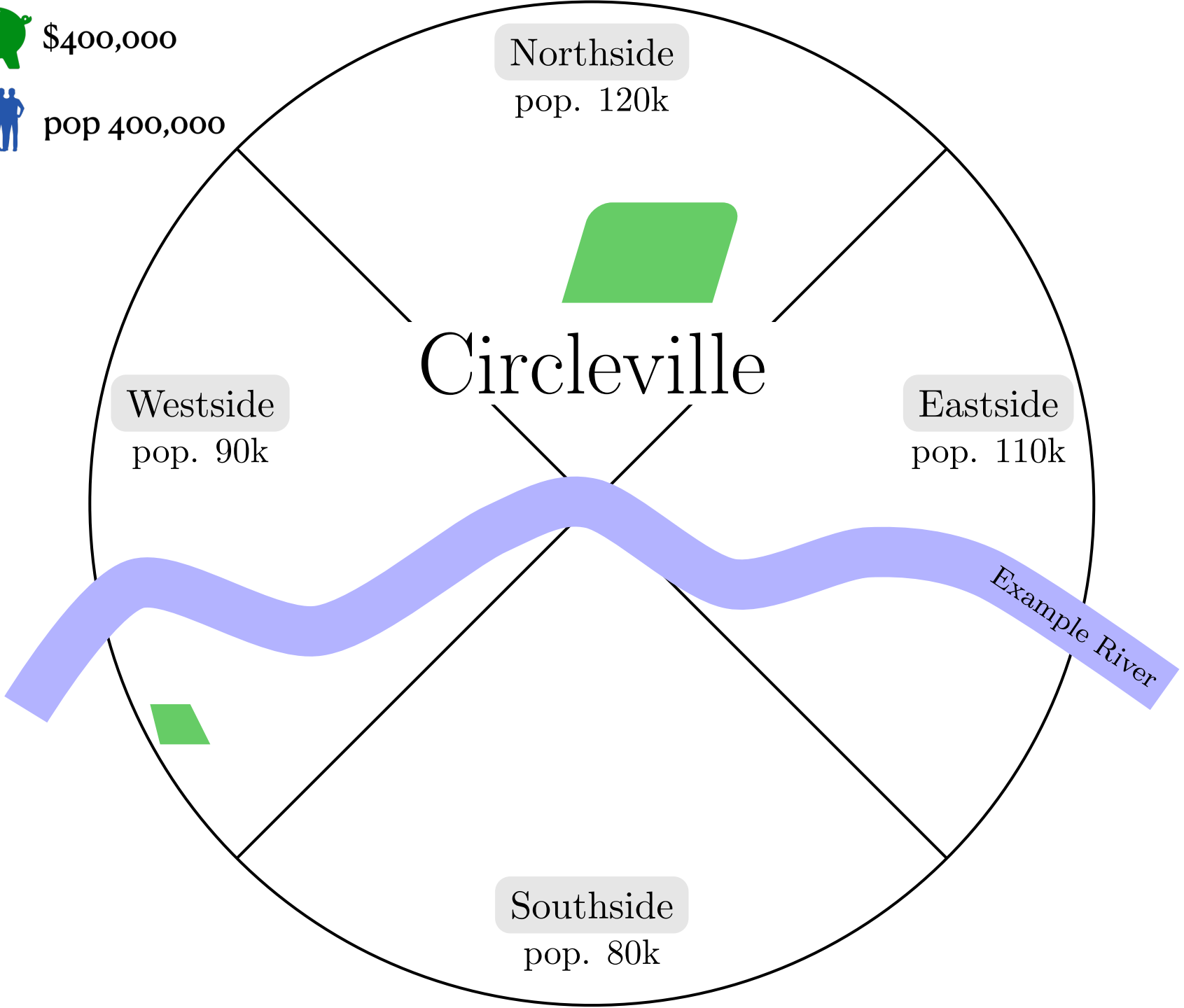






\$400,000

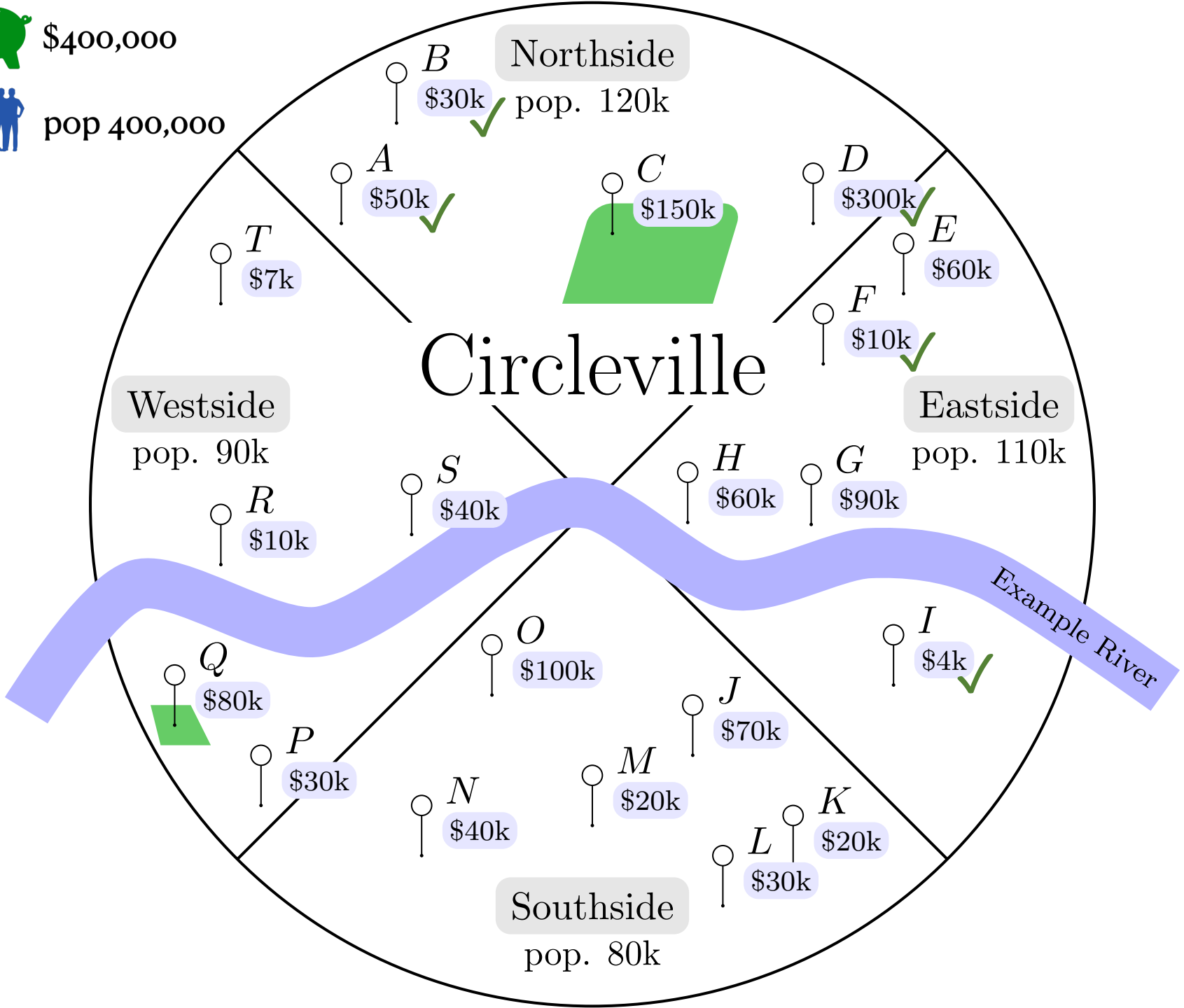


pop 400,000



 \$400,000

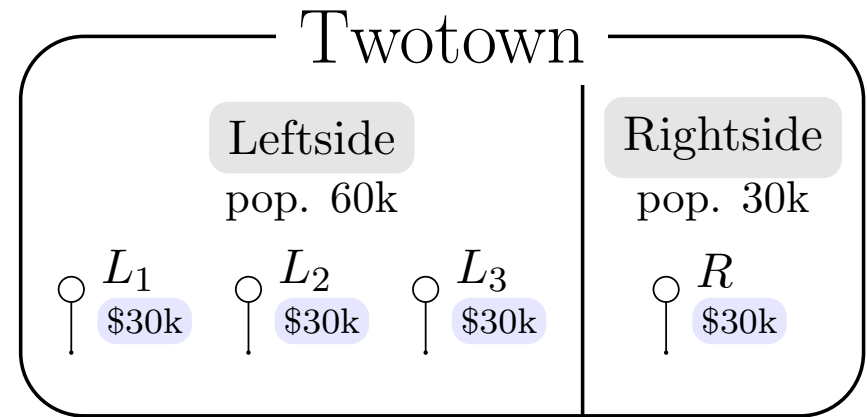
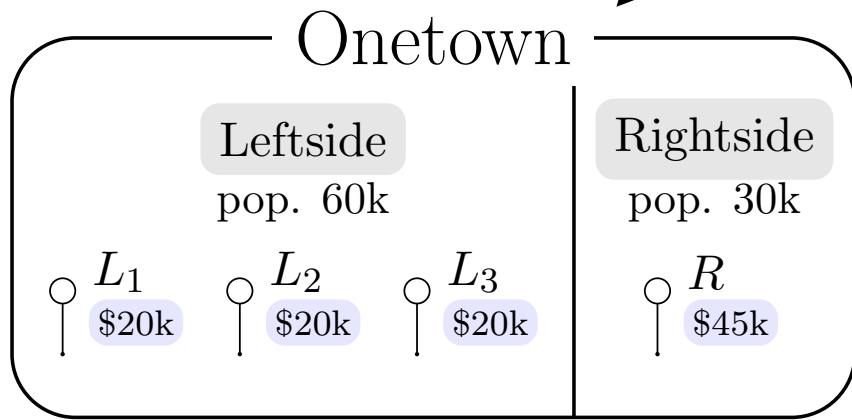
 pop 400,000



 \$90,000

 pop 90,000

same budget
same population
same district structure
same utilities
same projects
same feasible sets



$\{L_1, L_2, L_3\} \longrightarrow$ PAV-score 110,000

$\{L_1, L_2, R\} \longrightarrow$ PAV-score 120,000

 not proportional!
Leftside deserves \$60k

not proportional!

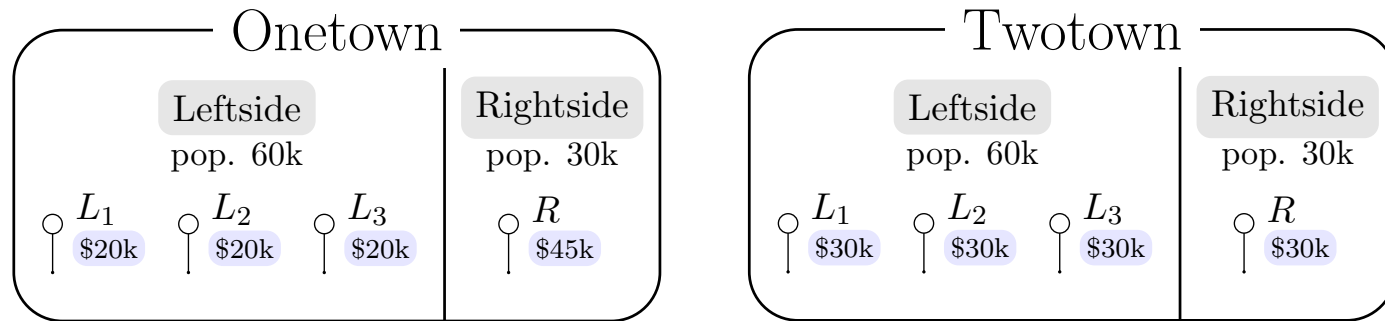
 Rightside deserves \$30k

$\{L_1, L_2, L_3\} \longrightarrow$ PAV-score 110,000

$\{L_1, L_2, R\} \longrightarrow$ PAV-score 120,000

Phragmén for PB

- Phragmén’s rule can easily be adapted:
 - Fill bank accounts
 - If the approvers of a project have enough money to **finance its cost**, implement the project
 - Stop when next project doesn’t fit into the budget.
- Picks correct outcome in Onetown and Twotown.



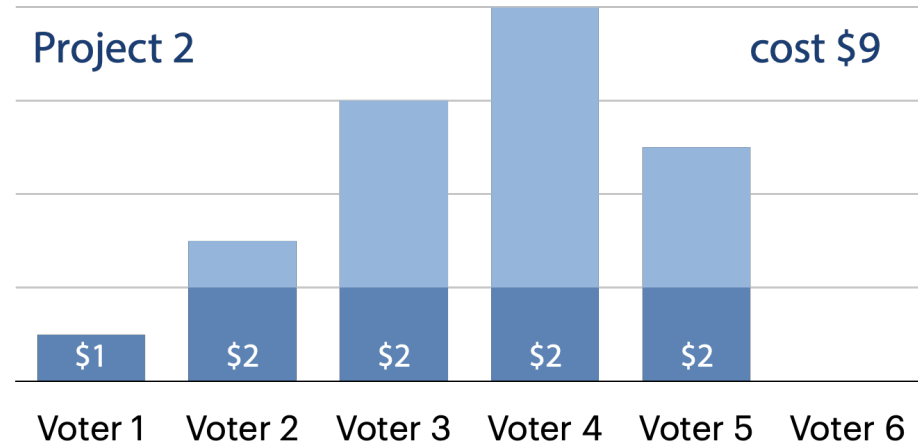
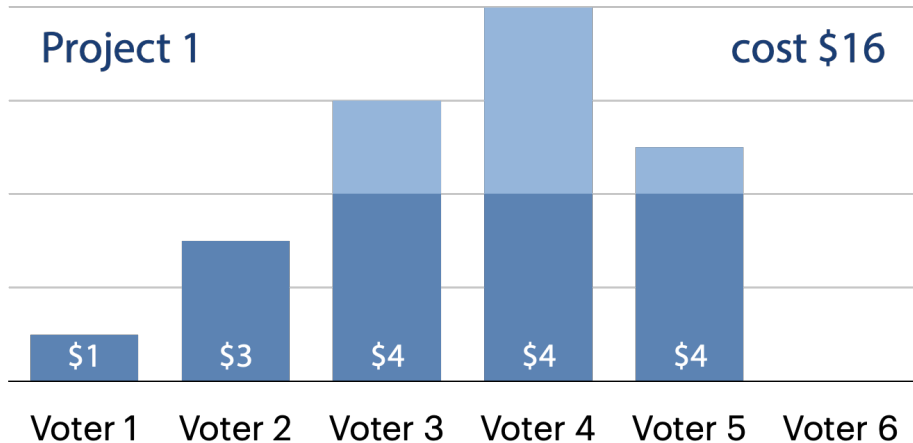
- **Satisfies “PJR”:** If $\frac{|S|}{n} \geq \alpha$ and $\text{cost}(\cap_{i \in S} A_i) \geq \alpha \cdot B$, then $\text{cost}((W \cap \cup_{i \in S} A_i) \cup \{c\}) \geq \alpha \cdot B$ for some $c \in \cap_{i \in S} A_i$.

Rule X for PB

- Split the city budget evenly among residents.
- Put each resident's share $\frac{B}{n}$ in a bank account.
- Repeatedly, until the budget runs out:
 - identify a project whose supporters have enough money left to afford it
 - charge the cost to supporters

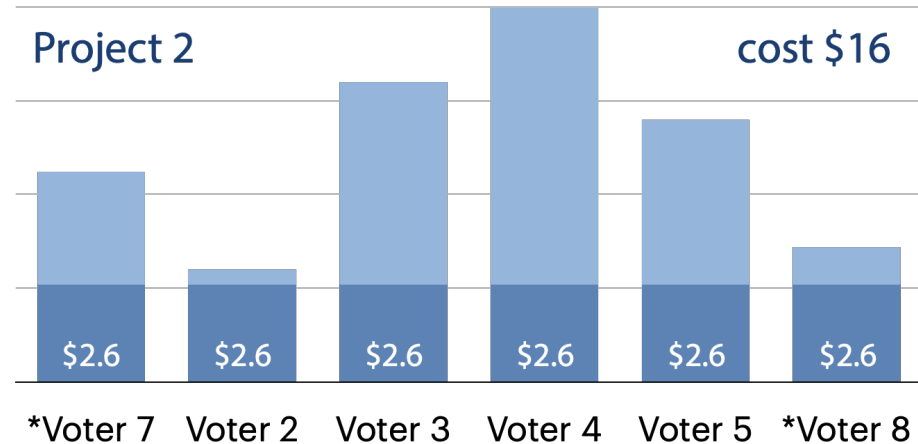
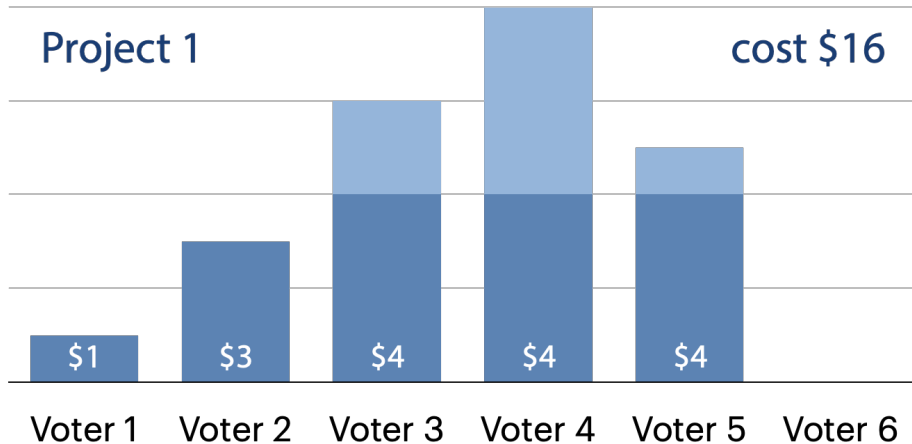
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 - always divide the cost of a project among supporters **as evenly as possible**
 - find an affordable project with the lowest max payment.



Rule X for PB

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Rule X for PB

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- Put each resident's share $\frac{B}{n}$ in a bank account.
- Repeatedly, until the budget runs out:
 - always divide the cost of a project among supporters **as evenly as possible**
 - find an affordable project with the lowest max payment.
- Rule X satisfies EJR!
 - Let $\frac{|S|}{n} \geq \alpha$, and take $T \subseteq \bigcap_{i \in S} A_i$ with $\text{cost}(T) \leq \alpha \cdot B$.
 - Then $u_i(W) \geq u_i(T)$ for some $i \in N$ (i.e. $|W \cap A_i| \geq |T|$.)

Additive Valuations

Motivating Example:

2019, Paris, 16th arrondissement

€560k: refurbish sports facility — 775 votes

€3k: materials for classroom project — 670 votes
— 1.15x as popular, 186x the cost!

- Utility of outcome: $u_i(W) = \sum_{c \in W} v_i(c)$.
- Phragmén: no obvious way of extending to additive utilities.
- Rule X: can extend using following idea: a voter's payment for a candidate should be proportional to the voter's utility for the candidate.
- Core may be empty!

Core for Additive Valuations

- A group $S \subseteq N$ with $\frac{|S|}{n} \geq \alpha$ *blocks* W if there is $T \subseteq C$ with $|T| \leq \alpha \cdot B$ such that $u_i(T) > u_i(W)$ for all $i \in S$.

	Voter 1	Voter 2	Voter 3	cost
$u_i(a)$	2	1	0	\$2
$u_i(b)$	0	2	1	\$2
$u_i(c)$	1	0	2	\$2
Budget B				\$3

- An approximation exists if we put $|T| \leq \frac{\alpha}{32} \cdot B$. The factor of 32 might be improvable to 2, but not further.
- Existence open for approval utilities.

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