

Optimized Democracy

Spring 2021 | Lecture 11

Rent Division

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ONCE UPON A TIME IN JERUSALEM



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PROVABLY FAIR SOLUTIONS.

Spliddit offers quick, free solutions to everyday fair division problems, using methods that provide indisputable fairness guarantees and build on decades of research in economics, mathematics, and computer science.



Share Rent



Split Fare



Assign Credit



Divide Goods



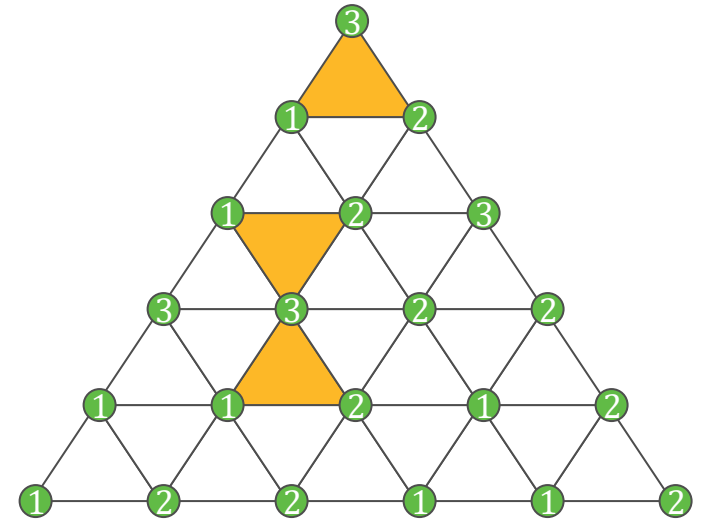
Distribute Tasks



Suggest an App

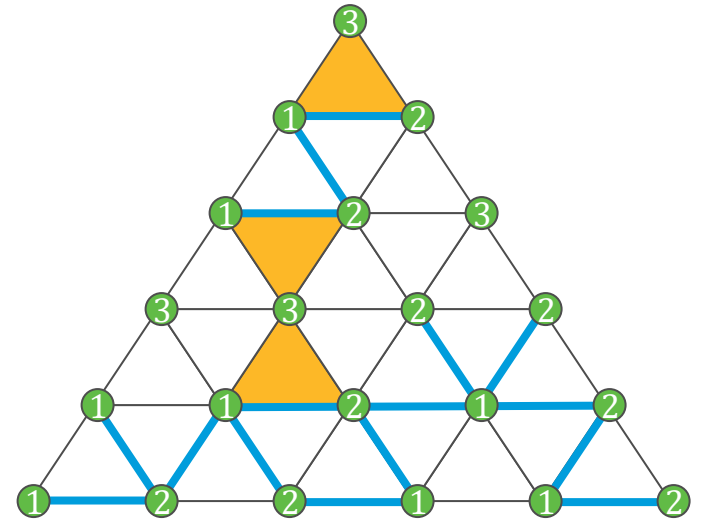
SPERNER'S LEMMA

- Triangle T partitioned into **elementary** triangles
- Label vertices by $\{1,2,3\}$ using **Sperner labeling**:
 - Main vertices are different
 - Label of vertex on an edge (i,j) of T is i or j
- **Lemma:** Any Sperner labeling contains at least one fully labeled elementary triangle



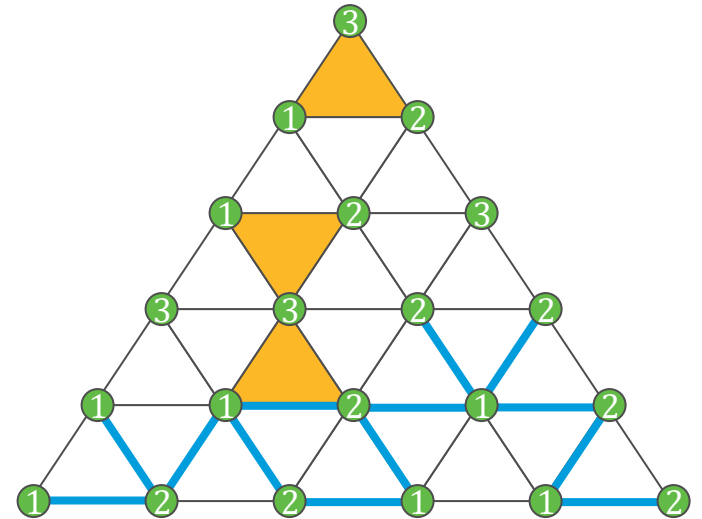
PROOF OF LEMMA

- Doors are 12 edges
- Rooms are elementary triangles
- #doors on the boundary of T is odd
- Every room has ≤ 2 doors; one door iff the room is 123



PROOF OF LEMMA

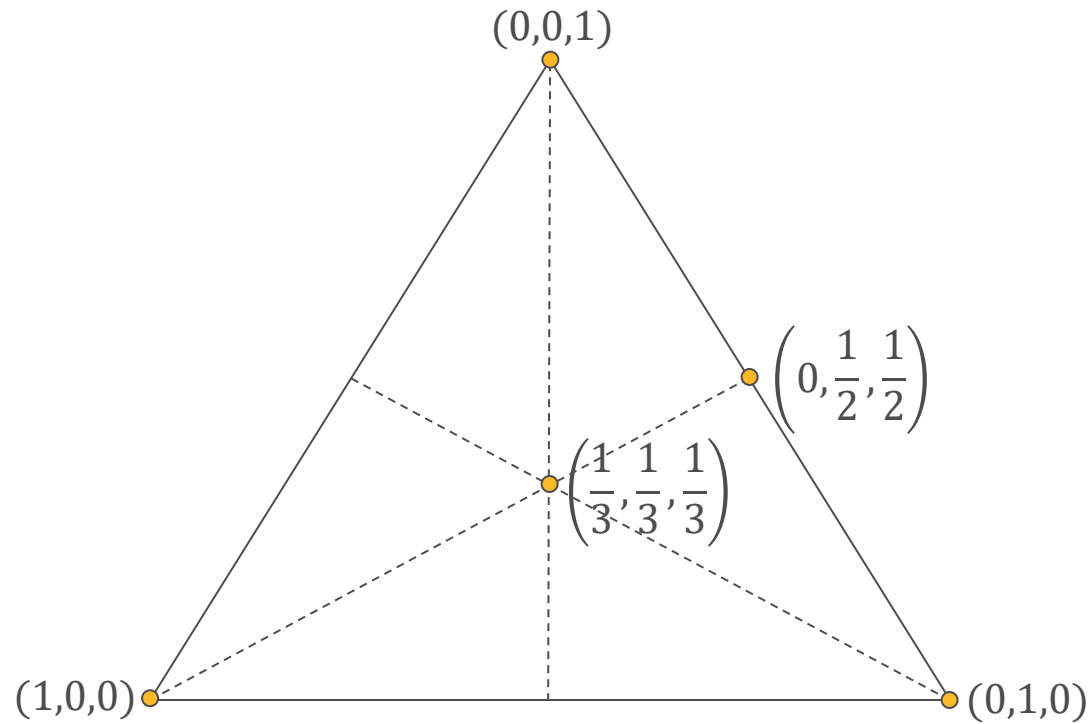
- Start at door on boundary and walk through it
- Room is fully labeled or it has another door...
- No room visited twice
- Eventually walk into fully labeled room or back to boundary
- But #doors on boundary is odd ■



THE MODEL

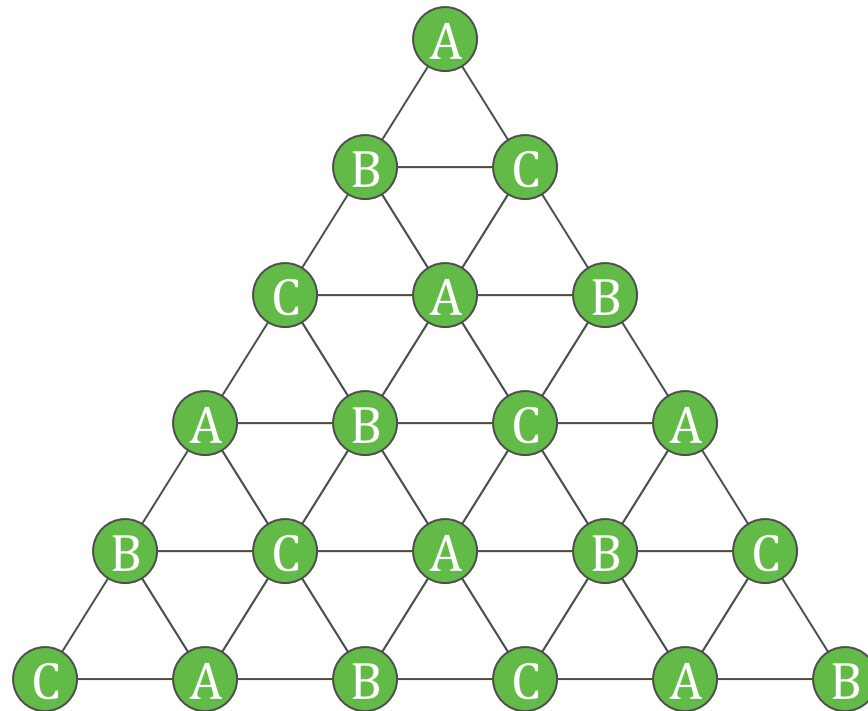
- Assume there are three players A, B, C
- Goal is to assign the rooms and divide the rent in a way that is **envy free**: each player prefers their own room at the given prices
- Sum of prices for three rooms is 1
- **Theorem**: An envy-free solution always exists under some assumptions

PROOF OF THEOREM



PROOF OF THEOREM

- “Triangulate” and assign “ownership” of each vertex to each of A, B, and C, in a way that each elementary triangle is an ABC triangle

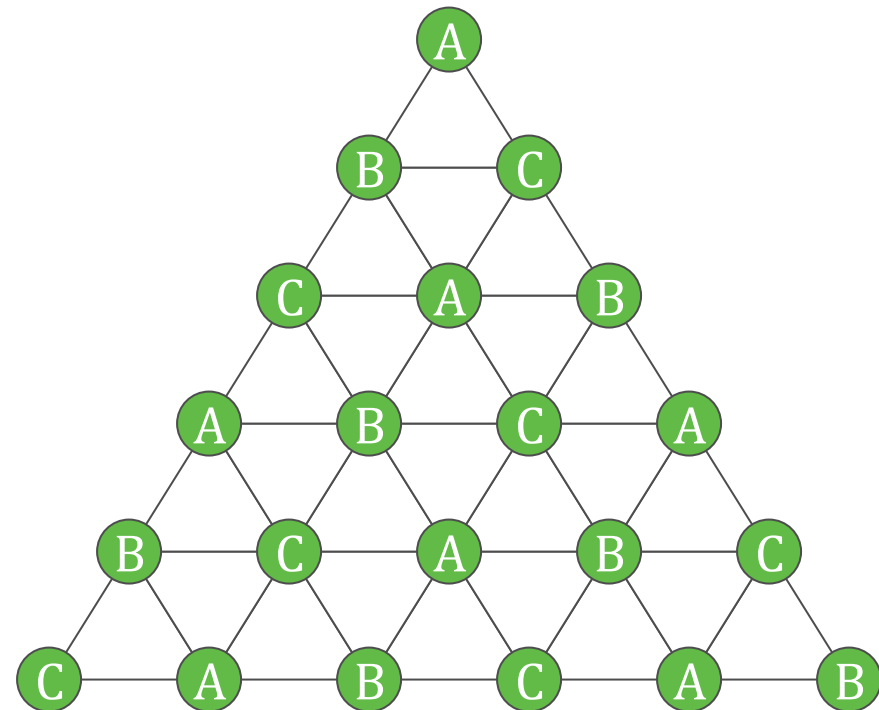
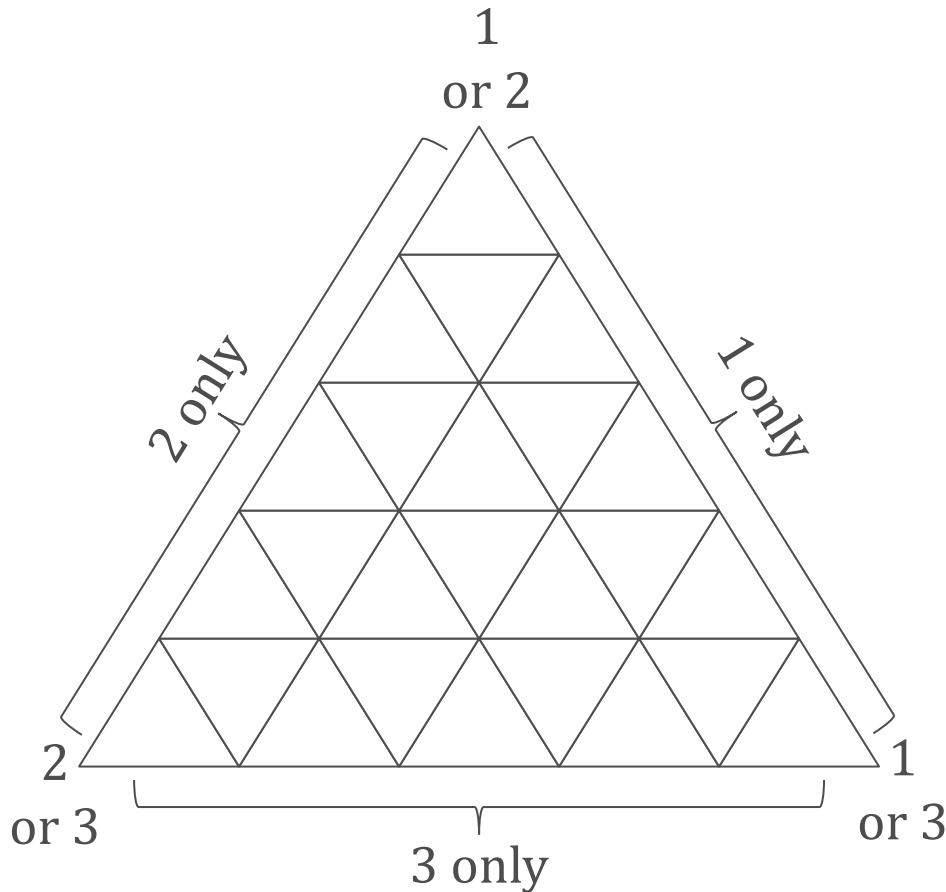


PROOF OF THEOREM

- Ask the owner of each vertex to tell us which room they prefer
- This gives a new labeling by 1, 2, 3
- **Assume** that a person wants a free room if one is offered to them

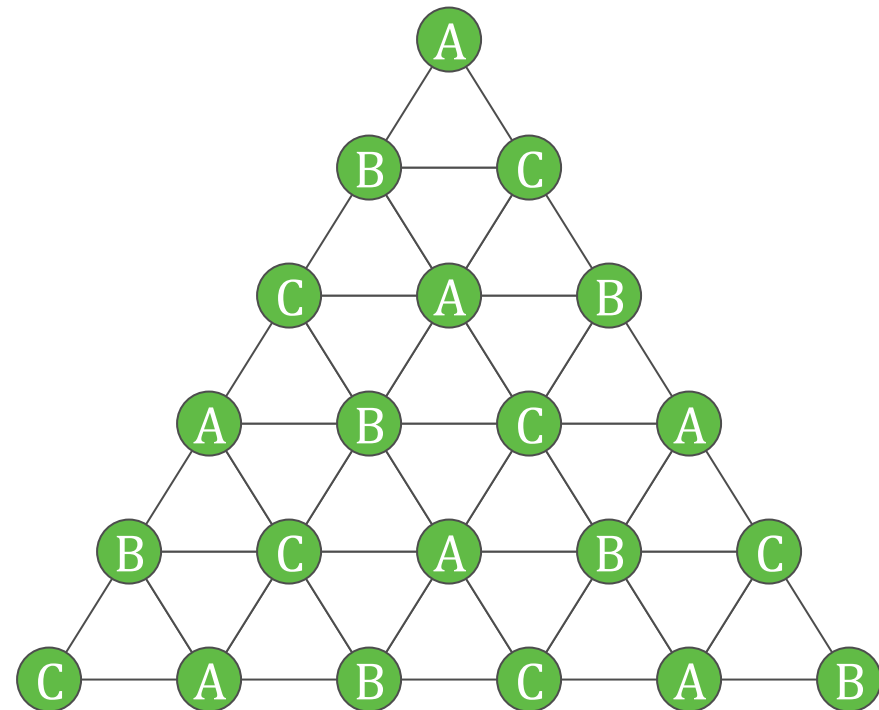
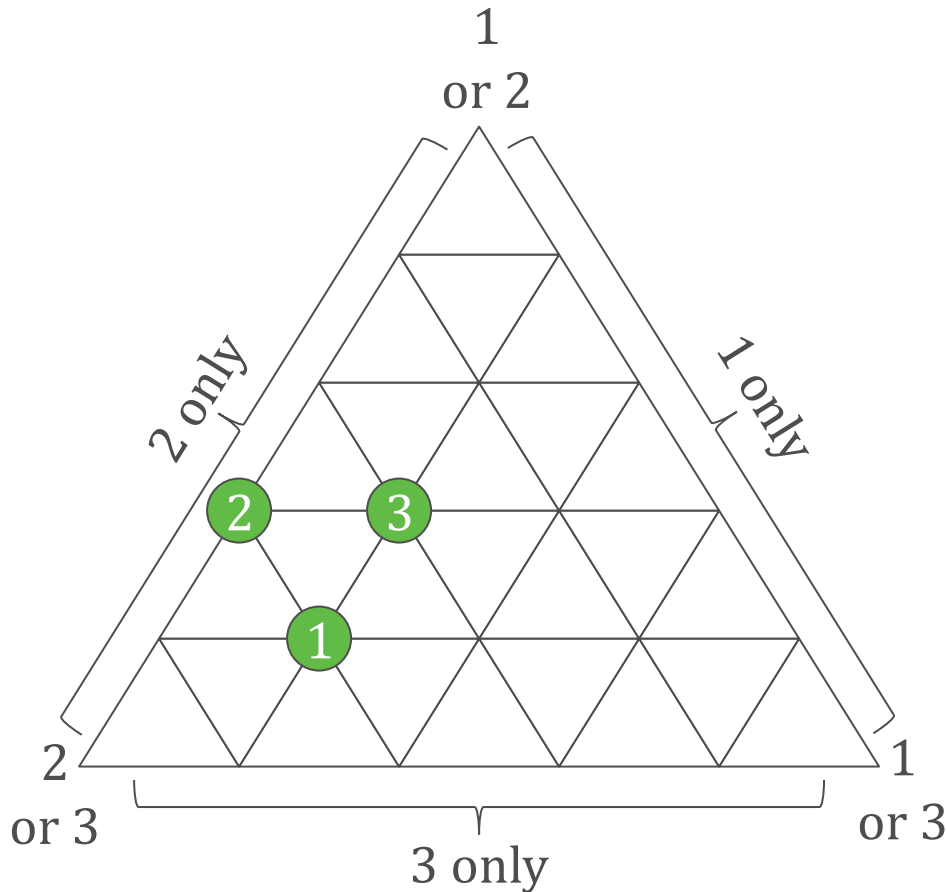
PROOF OF THEOREM

- Choice of rooms on edges is constrained by free room assumption



PROOF OF THEOREM

- Sperner's lemma (variant): such a labeling must have a 123 triangle



PROOF OF THEOREM

- Such a triangle is nothing but an approximately EF solution!
- By making the triangulation finer, we can approach envy-freeness
- Under additional closedness assumption, leads to existence of an EF solution ■

DISCUSSION

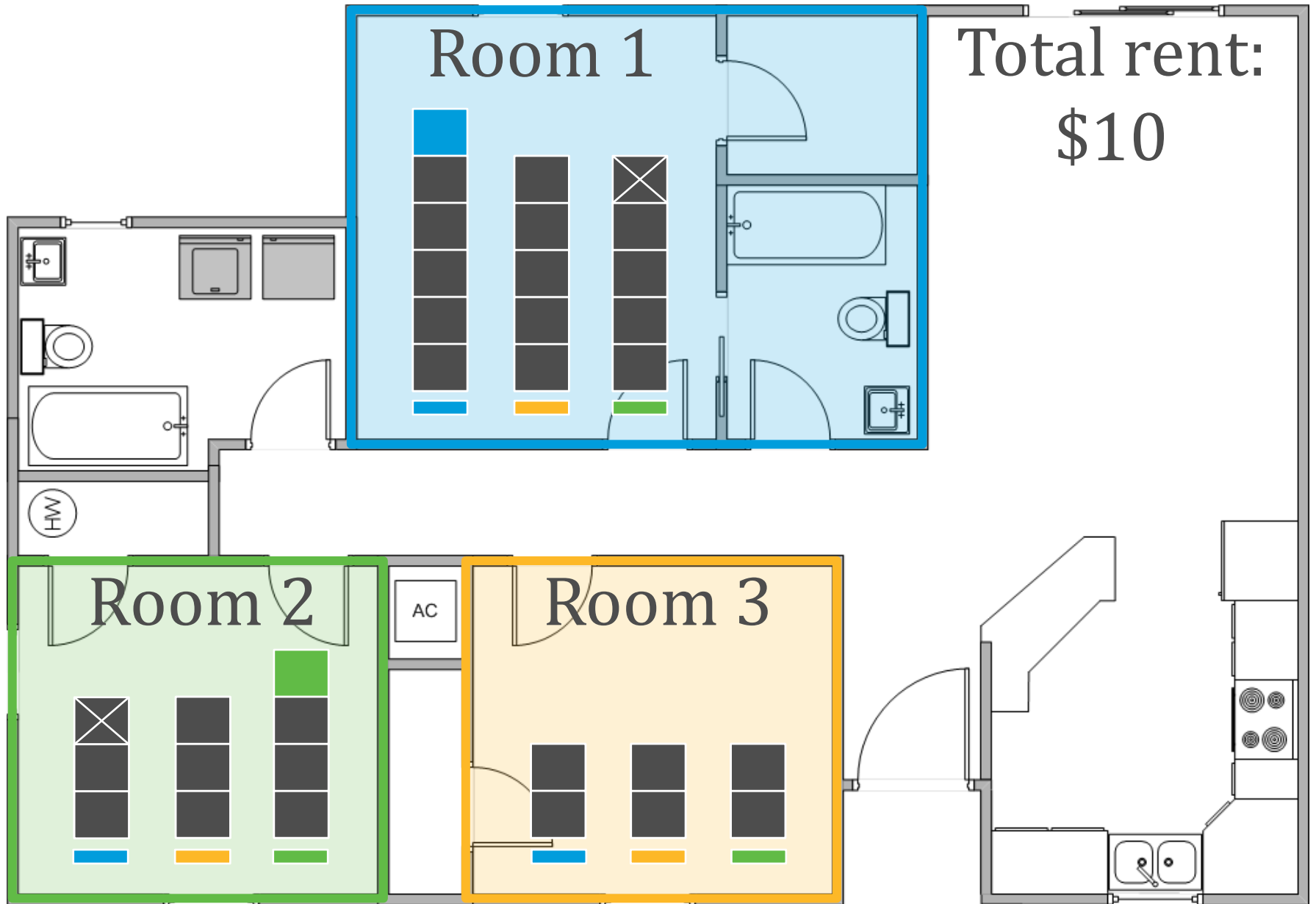
- It is possible to derive an algorithm from the proof
- Same techniques generalize to more players
- Same proof (with the original Sperner's Lemma) shows existence of EF cake division!

QUASI-LINEAR UTILITIES

- Suppose each player $i \in N$ has value v_{ir} for room r
- For all $i \in N$, $\sum_r v_{ir} = R$, where R is the total rent
- The utility of player i for getting room r at price p_r is $v_{ir} - p_r$
- A **solution** consists of an **assignment** π and a price vector \mathbf{p} , where p_r is the price of room r
- Solution (π, \mathbf{p}) is **envy free** if and only if
$$\forall i, j \in N, v_{i\pi(i)} - p_{\pi(i)} \geq v_{i\pi(j)} - p_{\pi(j)}$$
- **Theorem:** An envy-free solution always exists under quasi-linear utilities

Room 1

Total rent:
\$10



PROPERTIES OF EF SOLUTIONS

- Allocation π is **welfare-maximizing** if

$$\pi \in \operatorname{argmax}_{\sigma} \sum_{i \in N} v_{i\sigma(i)}$$

- Lemma 1:** If (π, \mathbf{p}) is an EF solution, then π is a welfare-maximizing assignment
- Lemma 2:** If (π, \mathbf{p}) is an EF solution and σ is a welfare-maximizing assignment, then (σ, \mathbf{p}) is an EF solution

PROOF OF LEMMA 1

- Let (π, \mathbf{p}) be an EF solution, and let σ be another assignment

- Due to EF, for all i ,

$$v_{i\pi(i)} - p_{\pi(i)} \geq v_{i\sigma(i)} - p_{\sigma(i)}$$

- Summing over all i ,

$$\sum_{i \in N} v_{i\pi(i)} - \sum_{i \in N} p_{\pi(i)} \geq \sum_{i \in N} v_{i\sigma(i)} - \sum_{i \in N} p_{\sigma(i)}$$

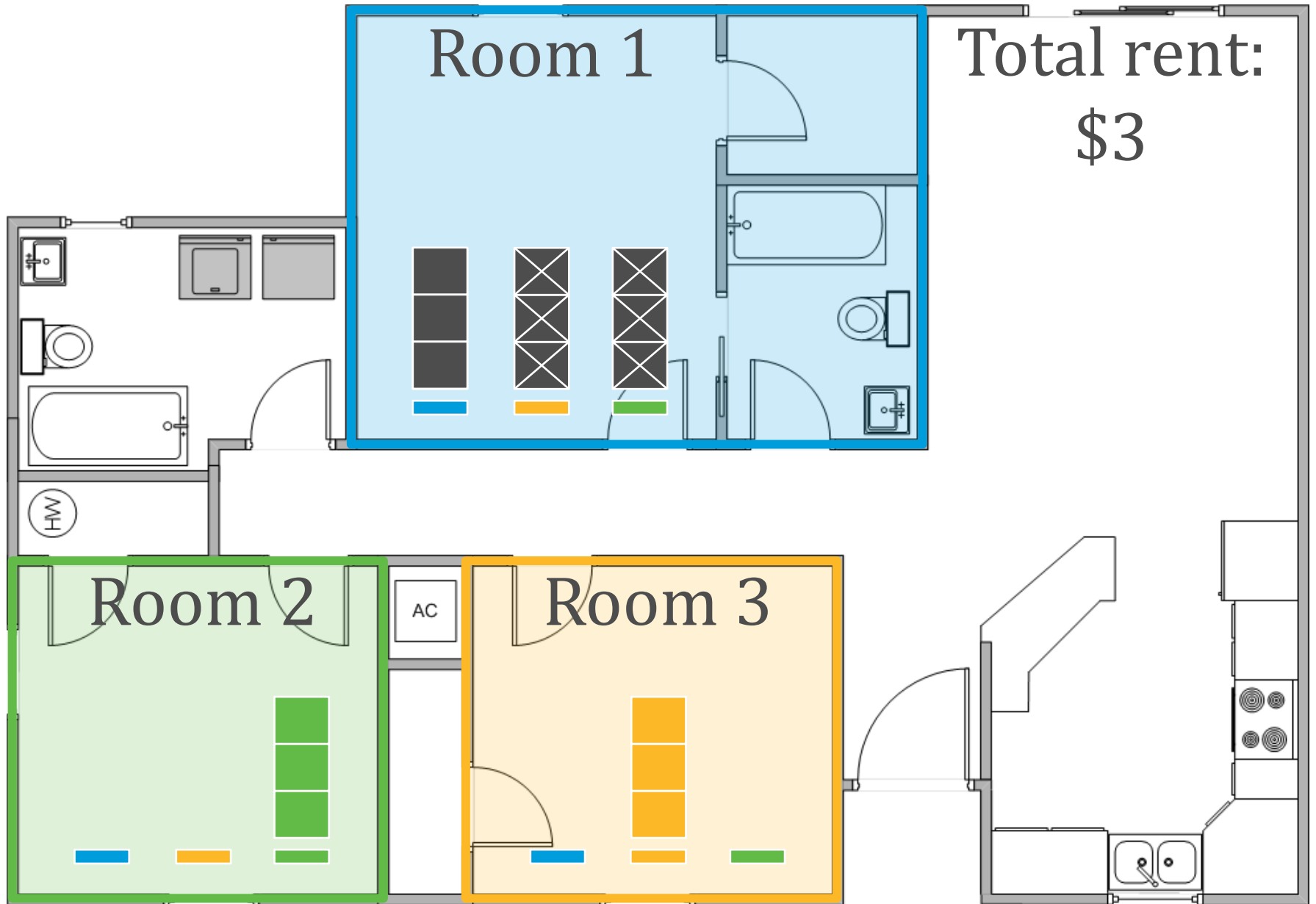
- We get the desired inequality because prices sum up to R ■

POLYNOMIAL-TIME ALGORITHM

- Consider the algorithm that finds a welfare-maximizing assignment π , and then finds prices \mathbf{p} that satisfy the EF constraint
- **Theorem:** The algorithm always returns an EF solution, and can be implemented in polynomial time
- **Proof:**
 - We know that an EF solution (σ, \mathbf{p}) exists, by Lemma 2 (π, \mathbf{p}) is EF, so we would be able to find prices satisfying the EF constraint
 - The first part is max weight matching, the second part is a system of linear inequalities ■

Room 1

Total rent:
\$3



OPTIMAL EF SOLUTIONS



Straw Man Solution

Max sum of utilities
Subject to envy freeness



Maximin Solution

Max min utility
Subject to envy freeness



Equitable solution

Min max difference in utils
Subject to envy freeness

OPTIMAL EF SOLUTIONS

- **Theorem:** The maximin and equitable solutions can be computed in polynomial time
- **Theorem:** The maximin solution is unique
- **Theorem:** The maximin solution is equitable, but not vice versa

CAVEAT: STRATEGYPROOFNESS

- Lemma 1 tells us that any EF solution is welfare maximizing
- Therefore, any EF solution is Pareto efficient
- But there is no rent division algorithm that is both strategyproof and Pareto efficient [Green and Laffont 1979]
- However, strategic behavior is largely a nonissue in practice in the rent division domain

CAVEAT: NEGATIVE RENT

- Envy-freeness may require negative rent, as the following example shows:

$$\begin{pmatrix} 36 & 34 & 30 & 0 \\ 31 & 36 & 33 & 0 \\ 34 & 30 & 36 & 0 \\ 32 & 33 & 35 & 0 \end{pmatrix}$$

- Whatever player i gets room 4 must pay 0, and the prices of the other rooms must be exactly their values to prevent envy
- Easy to verify that i can't be any of the players

DISCUSSION

- The first model makes no assumptions on utilities other than players preferring free rooms
- The second model assumes quasi-linear utilities

Question

What are some advantages and disadvantages of each of the two models?



INTERFACES

Divide Your Rent Fairly

APRIL 28, 2014

When you're sharing an apartment with roommates, it can be a challenge to decide who takes which bedroom, and at what price. Sit down with your roommates and use the calculator below to find the fair division. [RELATED ARTICLE](#)

What's your total rent? \$ 1000

How many of you are there? 2 3 4 5 6 7 8

If the rooms have the following prices, which room would you choose?

Choices will not necessarily be in order and the same roommate may be asked to choose multiple times in a row. Each roommate keeps choosing until a fair division is found.

Roommate A	\$250 Room 1	\$750 Room 2
Roommate B	<input type="checkbox"/> \$188 Room 1	<input type="checkbox"/> \$813 Room 2

Past Choices

		Room 1	Room 2
All	Roommate B	\$125.00	\$875.00
Roommate A	Roommate B	\$250.00	\$750.00
Roommate B	Roommate B	\$500.00	\$500.00

THE BASICS

+

ALICE'S EVALUATIONS

-

Alice, use the sliders or textboxes to place values on each room. Think of these values as bids: you will never pay more than what you bid, and in most cases you will pay less. However, your values must sum to the total monthly rent: \$1000. You can use the *rescale* button to automatically adjust your values to add up to the rent.

Master Bedroom \$ 0 .00

Basement \$ 0 .00

2nd Floor \$ 0 .00

RESET

RESCALE

CONTINUE

Current Total: \$0

Target: \$1000

BOB'S EVALUATIONS

+

CLAIRE'S EVALUATIONS

+

RESULTS

+

NY TIMES (rental harmony)

<https://www.nytimes.com/interactive/2014/science/rent-division-calculator.html>

Spliddit (quasi-linear utilities)

<http://www.spliddit.org/apps/rent>

BIBLIOGRAPHY

F. E. Su. **Rental Harmony: Sperner's Lemma in Fair Division.** American Mathematical Monthly, 1999.

A. Alkan and G. Demange and D. Gale. **Allocation of Indivisible Goods and Criteria of Justice.** Econometrica, 1991.

Y. Gal and M. Mash and A. D. Procaccia and Y. Zick. **Which Is the Fairest (Rent Division) of Them All?** Journal of the ACM, 2017.