# CMU 15-781 Lecture 7: Planning III

Teacher: Gianni A. Di Caro

### PLANNING GRAPHS

- Graph-based data structure representing a polynomial-size/time approximation of the exponential search tree
- Can be used to automatically produce good heuristic estimates (e.g., for A\*)
- Can be used to search for a solution using the GRAPHPLAN algorithm

## PLANNING GRAPHS

- Leveled graph: vertices organized into levels/stages, with edges only between levels
- Two types of *vertices* on alternating levels:
  - Conditions
  - Operations
- Two types of *edges*:
  - Precondition: from condition to operation
  - Postcondition: from operation to condition



•  $S_0$  contains all the conditions that hold in initial state



Level  $S_0$  Level  $O_0$ Precondition

• Add operation to level  $O_i$  if its preconditions appear in level  $S_i$ 



- Add condition to level  $S_i$  if it is the postcondition of an operation (it is in ADD or DELETE lists) in level  $O_{i-1}$
- Keep a previous condition of no action negates it (*persistence*, *no-op* action)

#### CONDITIONS MONOTONICALLY INCREASE



#### #Conditions

(always carried forward by no-ops)

Slide based on Brafman



• Repeat ...

15781 Fall 2016: Lecture 7



Idea:  $S_i$  contains all conditions that could hold at stage *i* based on past actions;  $O_i$  contains all operations that could have their preconditions satisfied at time *i* 

*No ordering* among the operations is assumed at each stage, they could be executed in parallel

- $\rightarrow$  The level *j* at which a condition first appears is a (good) estimate of how difficult is to achieve that condition
- $\rightarrow$  Can optimistically estimate how many steps it takes to reach a goal g (or sub-goal  $g_i$ ) from the initial state

#### **OPERATIONS MONOTONICALLY INCREASE**



#### #Operations

(as a result of conditions monotonic increase, that keep previous preconditions hold and set new preconditions true)

Slide based on Brafman

## MUTUAL EXCLUSION LINKS

- As it is the graph would be too optimistic!
- The graph also *records conflicts* between actions or conditions: two operations or conditions are mutually exclusive (mutex) *if no valid plan can contain both at the same time*
- A bit more formally:
  - *Two operations* are mutex if their preconditions or postconditions are mutex (*inconsistent effects, competing needs, interference*)
  - *Two conditions* are mutex if one is the negation of the other, or all action pairs that achieve them are mutex (*inconsistent support*)



• "Have cake and eat cake too" problem

Initial state: *Have*(*Cake*)

Goal:  $Have(Cake) \wedge Eaten(Cake)$ 

Eat(Cake):PRECOND: EFFECT:

 $\begin{array}{l} Have(Cake) \\ \neg Have(Cake) \land Eaten(Cake) \end{array}$ 

Bake(Cake): PRECOND: EFFECT:

 $\neg Have(Cake)$ Have(Cake)

15781 Fall 2016: Lecture 7



15781 Fall 2016: Lecture 7



15781 Fall 2016: Lecture 7

# MUTEX CASES

- Inconsistent postconditions (two ops): one operation negates the effect of the other; *Eat(Cake)* and no-op *Have(Cake)*
- Interference (two ops): a postcondition of one operation negates a precondition of other; Eat(Cake) and no-op Have(Cake) (issue in parallel execution, the order should not matter but here it would)



Inconsistent Postconditions







15781 Fall 2016: Lecture 7

# MUTEX CASES

- Competing needs (two ops): a precondition of one operation is mutex with a precondition of the other because they are the negate of each other, like for *Bake(Cake)* and *Eat(Cake)*, or because they have inconsistent support
- Inconsistent support (two conditions): each possible pair of operations that achieve the two conditions is mutex, *Have(Cake)* and *Eaten(Cake)*, are mutex in S<sub>1</sub> but not in S<sub>2</sub> because they can be achieved by *Bake(Cake)* and *Eaten(Cake)*





Inconsistent Support

#### 15781 Fall 2016: Lecture 7

Initial state: *Have*(*Cake*)

Goal:  $Have(Cake) \wedge Eaten(Cake)$ Eat(Cake): PRECOND: Have(Cake) $\neg Have(Cake) \land Eaten(Cake)$ EFFECT: Inconsistent support Competing Bake(Cake): needs **PRECOND:**  $\neg Have(Cake)$ Have(Cake)Effect: S<sub>0</sub> S<sub>2</sub> S, A<sub>0</sub> A, Bake(Cake) Have(Cake) Have(Cake Have(Cake) no-op no-op ≓no-op -Have(Cake) ¬Have(Cake) Eat(Cake) Eat(Cake) Eaten(Cake) Eaten(Cake) no-op -Eaten(Cake) no-op -Eaten(Cake ¬Eaten(Cake) no-op

15781 Fall 2016: Lecture 7

## DINNER DATE EXAMPLE

- Initial state: garbage  $\land$  cleanHands  $\land$  quiet
- Goals: dinner  $\land$  present  $\land \neg$ garbage
- Actions:
  - Cook: cleanHands  $\Rightarrow$  dinner
  - $\circ \quad \text{Wrap: quiet} \Rightarrow \text{present}$
  - Carry: none  $\Rightarrow \neg$  garbage  $\land \neg$  cleanHands
  - Dolly: none  $\Rightarrow \neg$  garbage  $\land \neg$  quiet

Slide based on Brafman

What's up

with the Wrap

precondition?



\* Slide based on Brafman

### DINNER DATE EXAMPLE\*



Slide based on Brafman

## A CONSTRUCTIVE PROCESS WITH NO CHOICES INVOLVED

- The process of constructing the planning graph does not require to choose among the actions
- $\rightarrow$  No combinatorial search is involved!
- At each O level, identify applicable actions, including the persistence ones
- At each S level identify all conditions that could result from the actions at the previous level
- Check for mutexes and add mutex link relations
- Check for level-off

#### # mutexes monotonically decrease

- Intuitive proof:
  - If two operations or conditions are a mutex at a given level, they will be also mutex for all previous levels (because of less conditions or actions to use to possibly avoid their conflicts)
  - At each level we should also think of all operations that cannot be executed and all actions that do not hold: these are in mutex with everything
  - By adding them over time the total # of mutexs decreases



#### #Conditions mutex

(new operations appear that can resolve inconsistent support conflicts)

Slide based on Brafman

## OPERATION MUTEXES MONOTONICALLY DECREASE

- Proof idea:
  - Different behavior for the different types of mutexes
    - Inconsistent postconditions and interference mutex are properties of the operations themselves  $\Rightarrow$  hold at every level
    - Competing needs depend on conditions at level S<sub>i</sub>: if two conditions are mutex because they are the negation of each other, then the mutex will hold at every level and so for the operations. Instead, if the conditions are mutex because of inconsistent support, by induction the monotonic increasing in #operations monotonically implies decrease in #mutexes



# LEVELING OFF

- As a corollary of the previous properties, we see that the planning graph levels off
  - Consecutive levels become identical

• Proof:

- $\circ$  Limit in the # of construction steps
- $\circ \quad \mbox{Upper bound on $\#$ operations and $$ $\#$ conditions based on problem definition $$ $$
- Lower bound of 0 on #mutexes

## PLANNING GRAPH COMPLEXITY

- The graph is polynomial in the size of the planning problem P(L,A)
  - Each  $S_i$  has at most L nodes and  $L^2$  mutex links
  - Each  $O_i$  has at most A+L nodes,  $(A+L)^2$  mutex links, and 2(AL+L) pre- and post-condition links
  - A graph with n levels has a size of  $O(n(A+L)^2)$
  - The time to build the graph has the same complexity

#### Planning graph $\cong$ Relaxed problem

- The relaxed problem provides an *optimistic estimate* of the number of steps to reach each sub-goal  $g_i \rightarrow$  Heuristics!
- Level cost of sub-goal (literal)  $g_i$  = level where  $g_i$  first appears
- Level cost is an *admissible* heuristic estimate but can be quite bad because level ≠ number of actions
- If any sub-goals  $g_i$  fail to appear in the final level of the graph  $\Rightarrow$  The problem is unsolvable ( $\infty$  LB)
- If all sub-goals appear  $\rightarrow$  "Possibly" there is a plan that achieves g, conditional that no mutexes exist

### HEURISTICS FROM PLANNING GRAPHS

- To estimate the cost of a conjunction of sub-goals:
  - Max level: max level cost of any goal (clearly admissible)
  - Level sum: sum of level costs for all goals
  - Set level: level at which all sub-goals appear without any *pair* of them being mutex
- Poll 3: Which is admissible:
  - 1. Level sum
  - 2. Set level
  - 3. Both
  - 4. Neither



## THE GRAPHPLAN ALGORITHM

- 1. Grow the planning graph until all sub-goals are reachable (appear as conditions) and *not mutex* (If planning graph levels off first, fail)
- 2. Call EXTRACT-SOLUTION on current planning graph
- 3. If none found, EXPAND-GRAPH by adding a level to the planning graph and try again

## THE GRAPHPLAN ALGORITHM

function GRAPHPLAN(problem) returns solution or failure  $graph \leftarrow INITIAL-PLANNING-GRAPH(problem)$   $goals \leftarrow CONJUNCTS(problem.GOAL)$ loop do

if goals all non-mutex in last level of graph then do
 solution ← EXTRACT-SOLUTION(graph,goals, NUMLEVELS(graph)
 if solution ≠ failure then return solution
 else if NO-SOLUTION-POSSIBLE(graph) then return failure
 graph ← EXPAND-GRAPH(graph,problem)

# EXTRACT-SOLUTION

- Backward search where each state corresponds to an S level of the planning graph and a set of unsatisfied goals
- Initial state is the last level of the planning graph, along with the goals of the planning problem
- Actions available at level  $S_i$  are to select any conflict-free subset of operations in  $A_{i-1}$  whose effects cover the goals in the state
- Resulting state has level  $S_{i-1}$  and its goals are the preconditions for selected actions
- Goal is to reach a state at level  $S_0$

#### FLAT TIRE PROBLEM

```
Init(At(Flat, Axle) \land At(Spare, Trunk))
Goal(At(Spare, Axle))
Action(Remove(Spare, Trunk),
   PRECOND: At(Spare, Trunk)
   EFFECT: \neg At(Spare, Trunk) \land At(Spare, Ground))
Action(Remove(Flat, Axle),
   PRECOND: At(Flat, Axle)
   EFFECT: \neg At(Flat, Axle) \land At(Flat, Ground))
Action(PutOn(Spare, Axle),
   PRECOND: At(Spare, Ground) \land \neg At(Flat, Axle)
   EFFECT: \neg At(Spare, Ground) \land At(Spare, Axle))
Action(LeaveOvernight,
   PRECOND:
   EFFECT: \neg At(Spare, Ground) \land \neg At(Spare, Axle) \land \neg At(Spare, Trunk)
            \wedge \neg At(Flat, Ground) \land \neg At(Flat, Axle))
```

#### BACKWARD SEARCH



15781 Fall 2016: Lecture 7

#### BACKWARD SEARCH, STEPS EXPLAINED

- Starting from  $S_0$ , that does not include the goal's literals, EXPAND-GRAPH is called twice. The literals from the goal are then present in the  $S_2$ , and EXPAND-SOLUTION is called
- Backward search starts at state  $S_2$  with goal At(Spare, Axle)
- The only action choice for achieving the goal is *PutOn(Spare, Axle)*
- The regression brings to a search state  $S_1$  with goals At(Spare, Ground),  $\neg At(Flat, Axle)$
- At(Spare, Ground) can only be achieved by Remove(Spare, Trunk)
- At(Flat,Axle) achieved either by Remove(Flat,Axle) or LeaveOvernight()
- LeaveOvernight() is mutex with  $Remove(Spare, Trunk) \rightarrow choose Remove(Spare, Trunk)$  and Remove(Flat, Axle)
- New search state at  $S_0$  with the goals At(Spare, Trunk), At(Flat, Axle), that are both in the initial conditions  $\rightarrow Solution found!$
- **Plan:** Remove(Spare, Trunk) and Remove(Flat, Axle) in level A<sub>0</sub>, PutOn(Spare, Axle) in A<sub>1</sub>.

### HEURISTICS FOR BACKWARD SEARCH

- How to choose among the actions?
- One approach: Greedy, based on the level cost of the literals in the set of sub-goals
  - Pick first the sub-goal with the highest level cost
  - Prefer actions with *easier preconditions* to achieve the selected sub-goal
  - $\circ$  Easier preconditions = Sum (or max) of the level costs of its preconditions is smallest

## GRAPHPLAN GUARANTEES

- Observation: The size of the *t*-level planning graph and the time to create it are polynomial in *t*, #operations, #conditions
- Theorem: GRAPHPLAN returns a plan if one exists, and returns failure if one does not exists

15781 Fall 2016: Lecture 7

# SUMMARY

- Terminology:
  - Planning graphs
  - $\circ$  Set level heuristic
- Algorithms:
  - GRAPHPLAN
  - FORWARD, BACKWARD SEARCH
- Big ideas:
  - Planning is search, but admits
    domain-independent heuristics

