

More Great Ideas in Theoretical CS

Learning Theory

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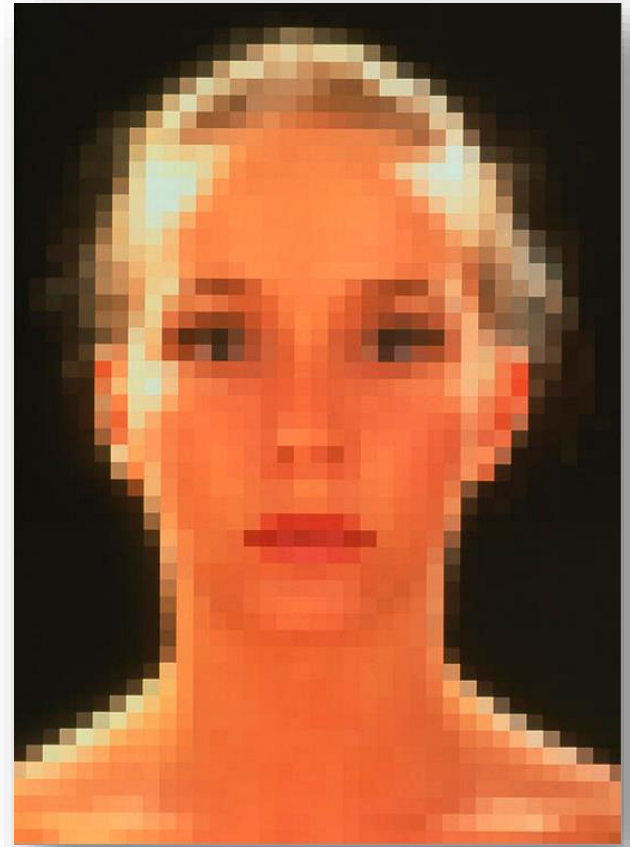
Ariel Procaccia (this time)

THE PAC MODEL

- Input space X
- D distribution over X : unknown but fixed
- Learner receives a set S of m instances x_1, \dots, x_m , independently sampled according to D
- Function class F of functions $f: X \rightarrow \{+, -\}$
- Assume target function $f_t \in F$
- Training examples $Z = \{(x_i, f_t(x_i))\}$

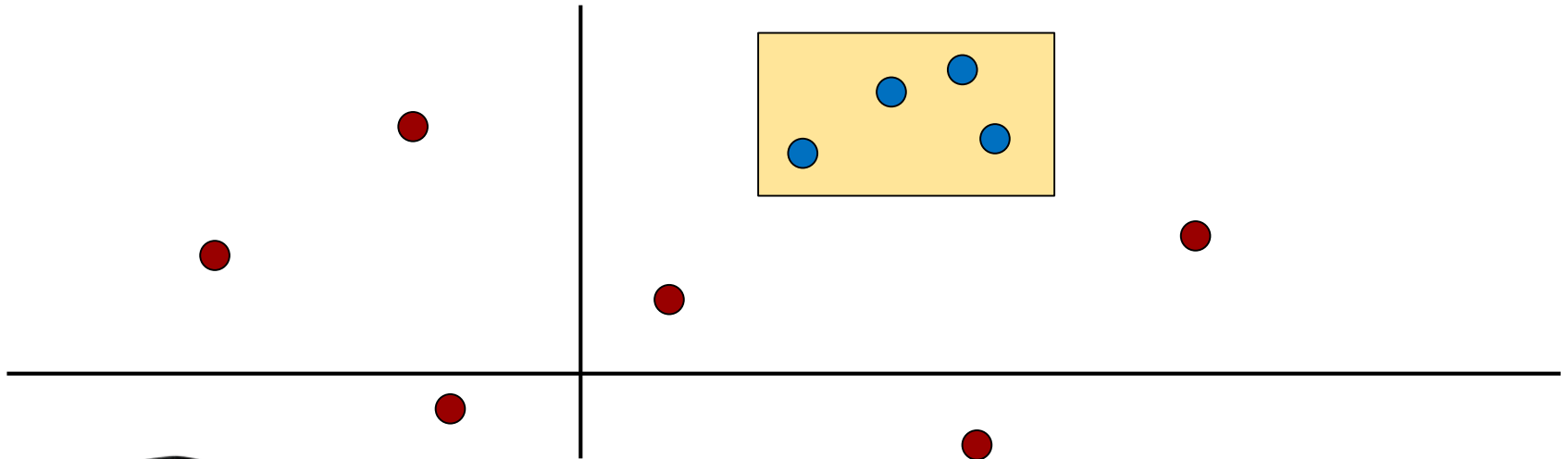
EXAMPLE: FACES

- $X = \mathbb{R}^{k \times \ell}$
- Each $x \in X$ is a matrix of colors, one per pixel
- $f_t(x) = +$ iff x is a picture of a face
- Training examples: Each is a picture labeled “face” or “not face”



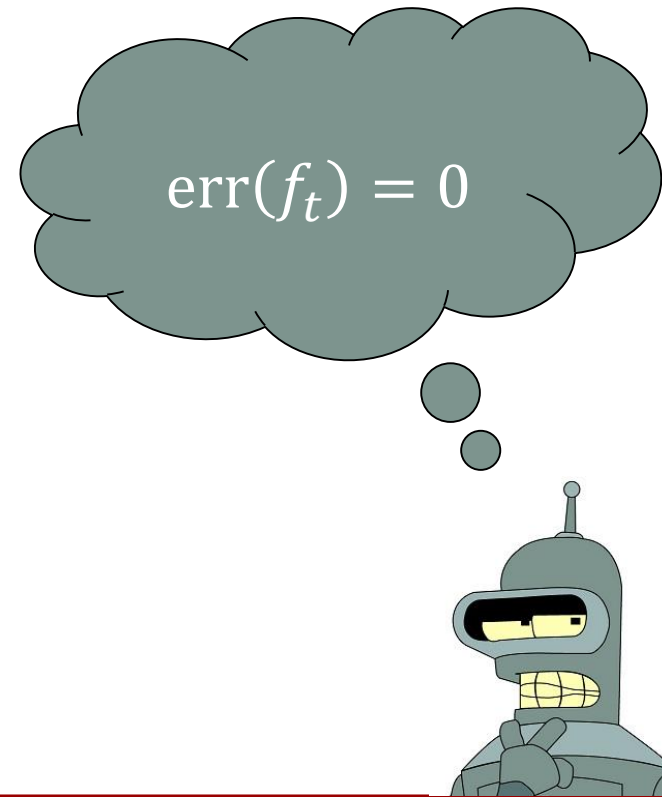
EXAMPLE: RECTANGLE LEARNING

- $X = \mathbb{R}^2$
- $F =$ axes-aligned rectangles
- $f(x) = +$ iff x is contained in f



THE PAC MODEL

- The **error** of function f is
$$\text{err}(f) = \Pr_{x \sim D} [f_t(x) \neq f(x)]$$
- Given **accuracy** parameter $\epsilon > 0$, would like to find function f with $\text{err}(f) \leq \epsilon$
- Given **confidence** parameter $\delta > 0$, would like to achieve $\Pr[\text{err}(f) \leq \epsilon] \geq 1 - \delta$



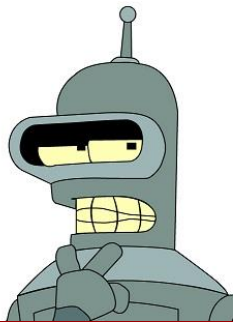
THE PAC MODEL

- A **learning algorithm** L is a function from training examples to F such that: for every $\epsilon, \delta > 0$ there exists $m^*(\epsilon, \delta)$ such that for every $m \geq m^*$ and every D , if m examples Z are drawn from D and $L(Z) = f$ then

$$\Pr[\text{err}(f) \leq \epsilon] \geq 1 - \delta$$

- F is **learnable** if there is a learning algorithm for F

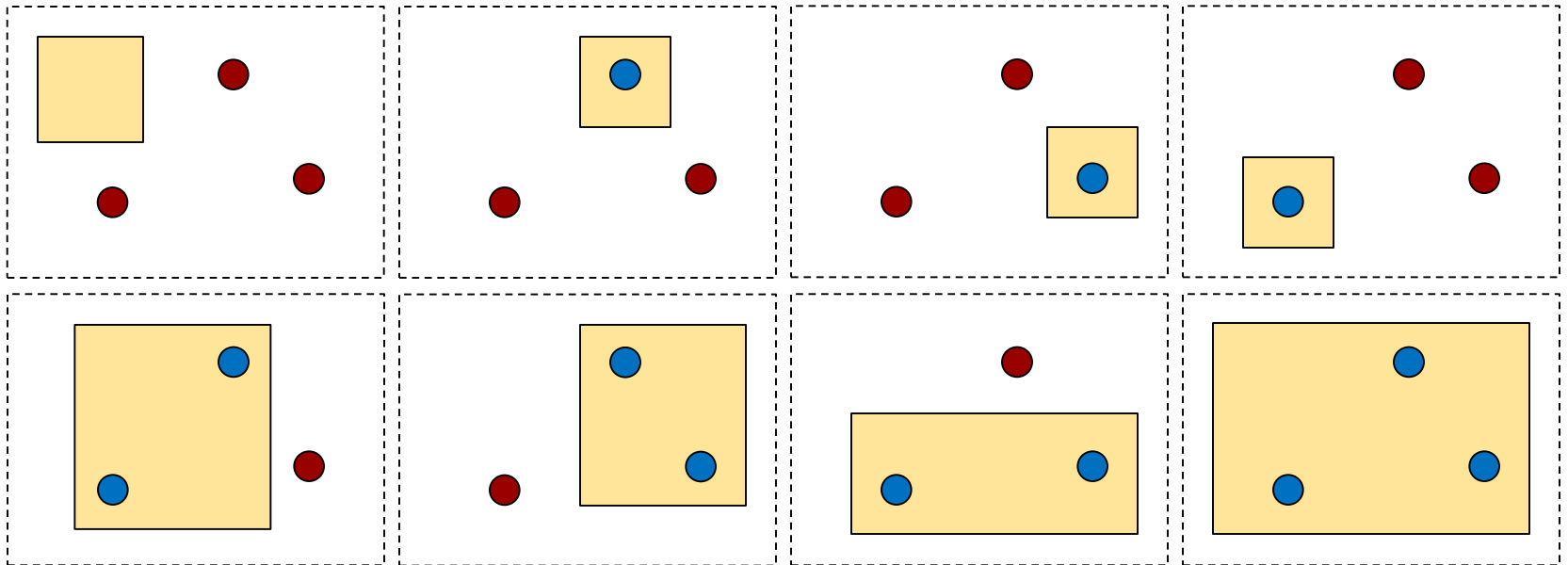
$m^*(\epsilon, \delta)$ is
independent of D !



VC DIMENSION


- We would like to obtain a general connection between learnability and combinatorial properties of the function class
- Let $S = \{x_1, \dots, x_m\}$
- $\Pi_F(S) = \{(f(x_1), \dots, f(x_m)) : f \in F\}$

VC DIMENSION



$$\Pi_F(S) = \{(-, -, -), (-, +, -), (-, -, +), (+, -, -), \\ (+, +, -), (-, +, +), (+, -, +), (+, +, +)\}$$

VC DIMENSION

- $X =$ real line
- $F =$ intervals; points inside interval are labeled by $+$, outside by $-$
- **Poll 1:** what is $|\Pi_F(S)|$ for $S =$ 

 - 1.
 - 2.
 - 3.
 - 4.

VC DIMENSION

• Poll 2: what is $|\Pi_F(S)|$ for $S =$ 

1. 5

2. 6

3. 7

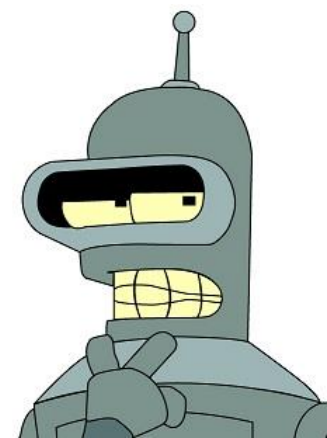
4. 8



VC DIMENSION

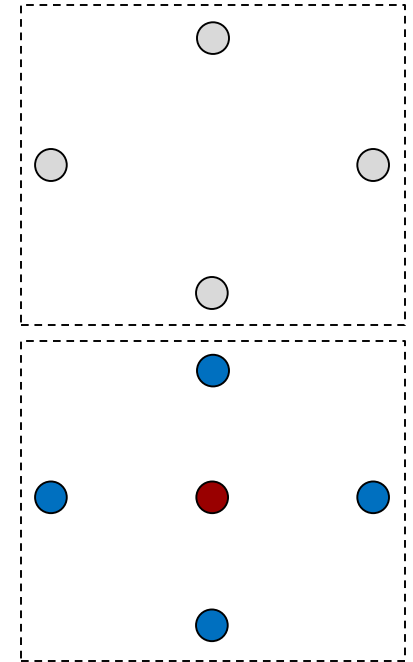
- S is **shattered** by F if $|\Pi_F(S)| = 2^{|S|}$
- The **VC dimension** of F is the cardinality of the largest set that is shattered by F

How do we
prove upper and
lower bounds?



EXAMPLE: RECTANGLES

- There is an example of four points that can be shattered
- For any choice of five points, one is “internal”
- A rectangle cannot label outer points by 1 and inner point by 0
- VC dimension is 4

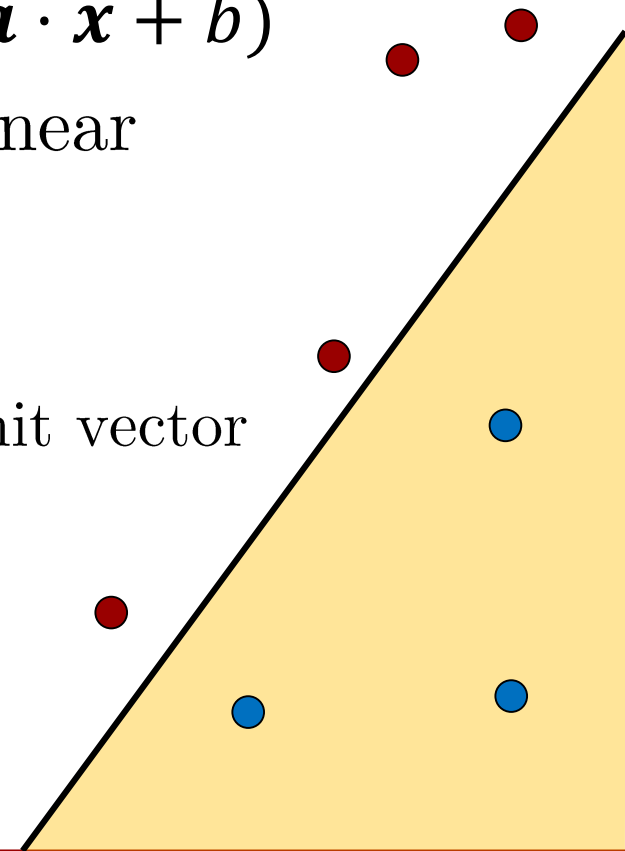


VC DIMENSION

- **Poll 3:** $X =$ real line, $F =$ intervals, what is $\text{VC-dim}(F)$?
 1. 1
 2. 2
 3. 3
 4. None of the above
- **Poll 4:** $X =$ real line, $F =$ unions of intervals, what is $\text{VC-dim}(F)$?
 1. 2
 2. 3
 3. 4
 4. None of the above

EXAMPLE: LINEAR SEPARATORS

- $X = \mathbb{R}^d$
- A linear separator is $f(\mathbf{x}) = \text{sgn}(\mathbf{a} \cdot \mathbf{x} + b)$
- **Theorem:** The VC dimension of linear separators is $d + 1$
- **Proof (lower bound):**
 - $\mathbf{e}^i = (0, \dots, 0, 1, 0, \dots, 0)$ is the i -th unit vector
 - $S = \{\mathbf{0}\} \cup \{\mathbf{e}^i : i = 1, \dots, d\}$
 - Given $y^0, \dots, y^d \in \{-1, 1\}$, set $\mathbf{a} = (y^1, \dots, y^d), b = y^0/2$ ■



SAMPLE COMPLEXITY

- If for any k there is a sample of size k that can be shattered by F , we say that $\text{VC-dim}(F) = \infty$
- **Theorem:** A function class F with $\text{VC-dim}(F) = \infty$ is not PAC learnable
- **Theorem:** Let F with $\text{VC-dim}(F) = d$. Let L be an algorithm that produces an $f \in F$ that is **consistent** with the given samples S . Then L is a learning algorithm for F with sample complexity

$$m^*(\epsilon, \delta) = O\left(\frac{1}{\epsilon} \log \frac{1}{\delta} + \frac{d}{\epsilon} \log \frac{1}{\epsilon}\right)$$