

## 15-251: Great Theoretical Ideas In Computer Science

### Recitation 13 Solutions

#### Euler is Even

Prove that for any  $n > 2$ ,  $\phi(n)$  is even.

Suppose  $k$  is relatively prime to  $n$ , so  $\gcd(k, n) = 1$ . We'll first show that  $\gcd(n - k, n) = 1$ .

Since  $\gcd(k, n) = 1$ , we know by Euler that there exist  $x, y$  so that  $kx + ny = 1$ . Thus we have  $1 = kx + ny = kx + ny - nx + nx = (k - n)x + n(x + y) = (n - k)(-x) + n(x + y)$ . Since  $x, y$  integers, so are  $x + y$  and  $-x$ .

Thus we have  $\gcd(n - k, n) \mid 1$ , so  $\gcd(n - k, n) = 1$ .

Now, we use this to pair off elements in  $Z_n^*$ .

First note that  $\frac{n}{2} \notin Z_n^*$ , since either  $n$  is odd or since  $n > 2$ ,  $\gcd(\frac{n}{2}, n) = \frac{n}{2} > 1$ .

Now, for each  $k \in Z_n^*$ ,  $n - k \in Z_n^*$ , and  $n - k \neq k$ . Thus  $\phi(n) = |Z_n^*|$  is even for  $n > 2$ .

#### RSA Practice

In lecture, we saw how RSA encryption is used. There are many important quantities used in this algorithm:

- $p, q$ : Two very large prime numbers.
- $n$ :  $n = pq$  is part of the public key
- $\phi(n)$ : Since  $p, q$  prime,  $\phi(n) = (p - 1)(q - 1)$
- $e$ :  $e$ , also part of the public key, is some member of  $Z_{\phi(n)}^*$
- $d$ :  $d$ , the private key, is the inverse of  $e$  in  $Z_{\phi(n)}^*$ , i.e.  $ed \cong_{\phi(n)} 1$
- $m$ : This is the message that will be sent

Let  $p = 17$ ,  $q = 7$ ,  $e = 11$

(a) Use the extended Euclidian Algorithm to find  $d$ .

First, we must find  $\phi(n)$ .  $\phi(n) = (p - 1)(q - 1) = 16 * 6 = 96$ . We must find  $d$  such that  $11d \equiv_{96} 1$ . We have:

$$\begin{aligned} 96 &= 11 \times 8 + 8 \\ 11 &= 8 \times 1 + 3 \\ 8 &= 3 \times 2 + 2 \\ 3 &= 2 \times 1 + 1 \end{aligned}$$

Now we work backwards:

$$\begin{aligned} 1 &= 3 - 2 \\ &= 3 - (8 - 3 \times 2) = 3 \times 3 - 8 \\ &= (11 - 8) \times 3 - 8 = (11 \times 3) - (8 \times 4) \\ &= (11 \times 3) - ((96 - 11 \times 8) \times 4) = (11 \times 35) - (96 \times 4) \end{aligned}$$

Thus,  $d = 35$ .

(b) Encrypt the message 3

First, we find  $n$  which is  $pq$  which is 119. We need to find  $3^{11} \pmod{119}$ . This is

$$\begin{aligned} 3^{11} &\equiv_{119} 3 * 3^5 * 3^5 \\ &\equiv_{119} 3 * 243 * 243 \\ &\equiv_{119} 3 * 5 * 5 \\ &\equiv_{119} 75 \end{aligned}$$

(c) Decrypt the message 2

We need to find  $2^{35} \pmod{119}$ . This is

$$\begin{aligned} 2^{35} &\equiv_{119} 2^{7^5} \\ &\equiv_{119} 128^5 \\ &\equiv_{119} 9^5 \\ &\equiv_{119} 3^5 * 3^5 \\ &\equiv_{119} 243 * 243 \\ &\equiv_{119} 5 * 5 \\ &\equiv_{119} 25 \end{aligned}$$

## Groups

Define  $\bullet : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$  as follows: For all  $x, y \in \mathbb{N}$ ,  
 $x \bullet 3 = x$

$$3 \bullet y = y$$

$$x \bullet y = x + y \text{ if both } x, y \neq 3$$

Is  $(\mathbb{N}, \bullet)$  a group?

For each of the four required properties of a group, prove or disprove that they hold for  $(\mathbb{N}, \bullet)$ .

Closure:

For all  $x, y \in \mathbb{N}$ ,  $x$ ,  $y$ , and  $x + y$  are also in  $\mathbb{N}$ . Thus closure is satisfied.

Associativity:

Not satisfied:  $(1 \bullet 2) \bullet 4 = 3 \bullet 4 = 4$ , but  $1 \bullet (2 \bullet 4) = 1 \bullet 6 = 7$ .

Identity:

This is satisfied, because 3 is an identity.  $\forall x \in \mathbb{N}$ ,  $x \bullet 3 = 3 \bullet x = x$ .

Inverses:

Not satisfied: Only 1, 2, and 3 have inverses. 0 has no inverse because  $0 \bullet x$  is  $x$  for  $x \neq 3$  and 0 if  $x = 3$ . Thus there is no  $x$  such that  $0 \bullet x = 3$ .

## Orders

Let  $G$  be an abelian group with operation  $\cdot$ .

Let  $x, y \in G$  have  $|x| = m$  and  $|y| = n$  with  $\gcd(m, n) = 1$ . Show that  $|x \cdot y| = mn$ .

We need to show that  $mn$  is the least  $k$  such that  $(x \cdot y)^k = e$ .

We have  $(x \cdot y)^{mn} = x^{mn} \cdot y^{mn}$  because  $G$  is abelian. Furthermore,  $x^{mn} = (x^m)^n = e^n = e$  and  $y^{mn} = (y^n)^m = e^m = e$ , so  $(x \cdot y)^{mn} = e$ .

Now, suppose  $(x \cdot y)^k = e$ . Then  $e = (x \cdot y)^{mk} = x^{mk} \cdot y^{mk} = y^{mk}$ , so  $n|mk$ . Similarly,  $e = (x \cdot y)^{nk} = x^{nk} \cdot y^{nk} = x^{nk}$ , so  $m|nk$ .

Thus  $mn|m^2k$  and  $mn|n^2k$ , so  $mn|\gcd(m^2k, n^2k) = k$ . Therefore  $mn \leq k$ , so  $|x \cdot y| = mn$ .